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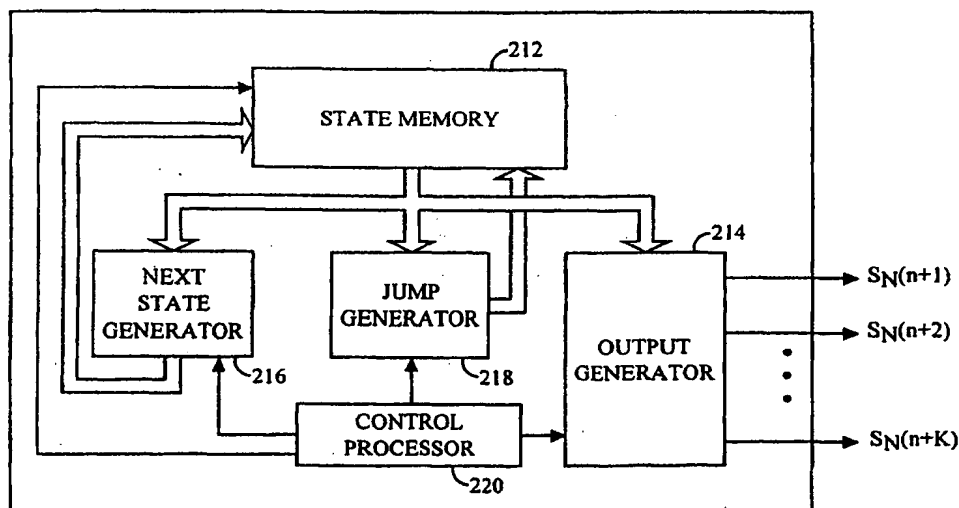
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EACH CLOCK PULSE BY COMPUTING THE BITS IN PARALLEL



(57) Abstract: A novel method and apparatus for generating PN sequences with an arbitrary number of bits, where the number of bits is provided in parallel with each clock pulse is described. This allows the sequences to be generated at high speed when needed, and allows parallel processing in the acquisition and demodulation processes. In the invention, the initial values of states are loaded into registers of a parallel PN generator, which immediately generates the next n bits of the PN sequence, where n is an arbitrary number dependent on performance required. Then, a first sub-part of the PN generator (406) of the present invention receives the present state of the PN generator (406) and outputs the state of the PN generator (406) n bits in the future.

A METHOD AND APPARATUS FOR GENERATING MULTIPLE BITS OF A PSEUDONOISE SEQUENCE WITH EACH CLOCK PULSE BY COMPUTING THE BITS IN PARALLEL

5 BACKGROUND OF THE INVENTION

I. Field of the Invention

The invention presented relates to pseudonoise (PN) sequence generators. More particularly, the present invention relates to a method and an
10 apparatus for generating PN sequence with each clock pulse by computing their bits in parallel.

II. Description of the Related Art

15 The Telecommunications Industry Association has standardized a method for code division multiple access (CDMA) communications in the IS-95 family of interim standards, entitled "Mobile Station-Base Station Compatibility Standard for Dual Mode Wideband Spread Spectrum Cellular System." In addition, the Telecommunications Industry Association in its submission to the
20 International Telecommunications Union, entitled "The cdma2000 ITU-R RTT Candidate Submission," describes proposed CDMA system that would be able to support higher data rates and higher capacity. Both in the IS-95 standard and in the cdma2000 proposal, the transmitted waveform is modulated in accordance with a pseudonoise spreading sequence.

25 The use of a pseudonoise sequence with appropriate autocorrelation characteristics is essential to the operation of a CDMA system in which multipath components are present. The generation and employment of pseudonoise sequences are described in detail in U.S. Patent No. 4,901,307, entitled "SPREAD SPECTRUM MULTIPLE ACCESS COMMUNICATION
30 SYSTEM USING SATELLITE OR TERRESTRIAL REPEATERS," assigned to the assignee of the present invention, and incorporated by reference herein. The use of CDMA techniques in a multiple access communication system is further disclosed in U.S. Patent No. 5,103,459, entitled "SYSTEM AND METHOD FOR GENERATING SIGNAL WAVEFORMS IN A CDMA CELLULAR
35 TELEPHONE SYSTEM," assigned to the assignee of the present invention, and incorporated by reference herein.

The aforementioned U.S. Patents Nos. 4,901,307 and 5,103,459 describe the use of a pilot signal used for acquisition. The use of a pilot signal enables

the remote user to acquire local base station communication system in a timely manner. The remote user gets synchronization information and relative signal power information from the received pilot signal. U.S. Patents Nos. 5,644,591 and 5,805,648, both entitled "METHOD AND APPARATUS FOR PERFORMING SEARCH ACQUISITION IN A CDMA COMMUNICATION SYSTEM," describe a novel and improved method and apparatus that reduces the remote user forward link acquisition time. Both patents are assigned to the assignee of the present invention and are incorporated by reference herein.

Space or path diversity is obtained by providing multiple signal paths through simultaneous links from a remote user through two or more cell-sites. Furthermore, path diversity may be obtained by exploiting the multipath environment through spread spectrum processing by allowing a signal arriving with different propagation delays to be received and processed separately. Examples of path diversity are illustrated in U.S. Patent No. 5,101,501, entitled "SOFT HANDOFF IN A CDMA CELLULAR TELEPHONE SYSTEM," and U.S. Patent No. 5,109,390, entitled "DIVERSITY RECEIVER IN A CDMA CELLULAR TELEPHONE SYSTEM," both assigned to the assignee of the present invention, and incorporated by reference herein.

In CDMA communications systems, a pilot signal is transmitted that allows a receiver to coherently demodulate the received signal. Within demodulator of such receivers is a channel estimate generator, which estimates the channel characteristics based on the pilot signal transmitted with values known to both the transmitter and the receiver. The pilot signal is demodulated and the phase ambiguities in the received signal are resolved by taking the dot product of the received signal and the pilot signal channel estimate. An exemplary embodiment of a circuit for performing the dot product operation is disclosed in U.S. Patent No. 5,506,865, entitled "PILOT CARRIER DOT PRODUCT CIRCUIT," assigned to the assignee of the present invention, and incorporated by reference herein.

SUMMARY OF THE INVENTION

The invention presented is a novel method and apparatus for generating a PN sequences with an arbitrary number of bits, where the number of bits is provided in parallel with each clock pulse. This allows the sequences to be generated at high speed when needed, and allows parallel processing in the acquisition and demodulation processes. The invention describes in detail generation of PN sequences as standardized for the IS-95 communications systems. As proposed in the IS-95 standards, the pseudonoise spreading

sequences are maximal length sequences that are capable of being generated using linear feedback shift-registers (LSFRs). Using a linear feedback shift-register, the PN sequences are computed one bit with each clock pulse.

5 In the invention, the initial PN states are loaded into registers of a parallel PN generator, which immediately generates the next n bits of the PN sequence, where n is an arbitrary number dependent on performance required. In addition, the present invention provides a method of determining the register states of the parallel PN generator an arbitrary number of cycles in the future. Thus, the present invention takes the present state of the registers of the
10 PN generator and outputs the next n bits of the generator. In addition, the PN generator of the present invention receives the present state of the PN generator and outputs the state of the PN generator n bits in the future. In this fashion, the entire PN sequence can be continuously generated.

15 It will be understood by one skilled in the art that although the present invention is directed toward the generation of a pseudonoise sequences compliant with systems standardized by the Telecommunications Industry Association, the teachings of the present invention are equally applicable to the generation of other pseudonoise sequences such as, the orthogonal Gold code sequences proposed for use in the W-CDMA, proposals to the International
20 Telecommunications Industry Association, proposals by the European Telecommunications Standards Institute (ETSI), and the Association of Radio Industries and Business (ARIB).

BRIEF DESCRIPTION OF THE DRAWINGS

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The features, objects, and advantages of the present invention will become more apparent from the detailed description set forth below when taken in conjunction with the drawings in which like reference characters identify correspondingly throughout and wherein:

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FIG. 1 illustrates a prior art embodiment of pseudonoise (PN) generators employing linear feedback shift-registers;

FIG. 2 depicts prior art of pseudonoise generators employed to generate parallel groups of PN sequence;

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FIG. 3 is a block diagram illustrating the generalized operation of the present invention apparatus for generating the PN sequences;

FIG. 4 shows one embodiment of the invention;

FIG. 5 is a simplified block diagram of an exemplary receiver chain using PN generators in accordance with the invention; and

FIG. 6 is a block diagram of a part of an exemplary single demodulation chain using PN generators in accordance with the invention.

DETAILED DESCRIPTION OF PREFERRED EMBODIMENTS

FIG. 1a illustrates a traditional apparatus employing a linear feedback shift-register for generating pseudonoise sequences. The generalized shift-register 100 from FIG. 1a comprises memory elements 102a, 102b, . . . , 102n, holding state values $S_0(n)$, $S_1(n)$, . . . , $S_N(n)$. The last value S_N constitutes an output of the shift-register, and also a feed-back to modulo-2 adders 104a, . . . , 104m. Before the value S_N is provided to a particular modulo-2 adder 104a, . . . , 104m, it is multiplied by an associated coefficient g_0, g_1, \dots, g_N . A coefficient will take a value of '1' if a feedback is desired, and a value of '0' otherwise.

Short-code pseudonoise sequences are used to modulate and demodulate the in-phase (I) and quadrature-phase (Q) components of the CDMA waveform. The I and Q short-code PN sequences are periodic with a period of $2^{15} - 1$ with a bit stuffed at the preamble of sequence to make the sequence periodic with an even factor of 2.

The short-code PN_I sequence satisfies a linear recursion specified by the following generator polynomial (P_I):

$$P_I(x) = x^{15} + x^{13} + x^9 + x^8 + x^7 + x^5 + 1. \quad (1)$$

FIG. 1.b depicts a shift-register implementation for generating the PN_I sequence. Note that in accordance with FIG. 1a, only the '1' valued coefficients $g_{15}, g_{13}, g_9, g_8, g_7, g_5, g_0$ are present.

The short-code PN_Q sequence satisfies a linear recursion specified by the following generator polynomial (P_Q):

$$P_Q(x) = x^{15} + x^{12} + x^{11} + x^{10} + x^6 + x^5 + x^4 + x^3 + 1. \quad (2)$$

FIG. 1.c depicts a shift-register implementation for generating the PN_Q sequence.

FIG. 1c shows a shift-register implementation of a long-code PN generator with a mask. The long-code is periodic, with period $2^{42} - 1$ chips and satisfies a linear recursion specified by the following characteristic polynomial (P):

$$P(x) = x^{42} + x^{35} + x^{33} + x^{31} + x^{27} + x^{26} + x^{25} + x^{22} + x^{21} + x^{19} + x^{18} + x^{17} + x^{16} + x^{10} + x^7 + x^6 + x^5 + x^3 + x^2 + x + 1 \quad (3)$$

5 The mask used for the long-code is channel type dependent, and can be found along with further details about the implementation of the PN generators in a document entitled "Physical Layer Standard for cdma2000 Spread Spectrum Systems."

10 It is sometimes desired to obtain an output of a shift-register as a parallel combination of output state values $S_N(n)$, $S_N(n+1)$, \dots , $S_N(n+K)$. FIG. 2 shows a block diagram of a parallel PN generator 200 according to the prior art. The PN generator comprises a shift-register 100 in accordance with a description for FIG. 1a, followed by a serial-to-parallel converter 202. The PN generator
15 outputs K values of $S_N(n)$ for shift instances n , $n+1$, \dots , $n+K$. However, there are K clock cycles required for generating the set of K output values. In the prior art understanding, in order to generate the parallel PN generator outputs, the outputs of the linear feedback shift-registers illustrated in FIGS. 1a and 1b are provided to the serial to parallel converter.

20

FIG. 3 shows a block diagram of inventive alternative to the implementation of FIG. 2. In general, a relationship between values of shift register in a state (n) and next state ($n+1$) can be expressed as a system of equations:

25

$$S_N(n+1) = g_{11} \cdot S_N(n) + \dots + g_{1N-1} \cdot S_2(n) + g_{1N} \cdot S_1(n) \quad (4a)$$

$$S_2(n+1) = g_{N-11} \cdot S_N(n) + \dots + g_{N-1N-1} \cdot S_2(n) + g_{2N} \cdot S_1(n) \quad (4n-1)$$

30

$$S_1(n+1) = g_{N1} \cdot S_N(n) + \dots + g_{NN-1} \cdot S_{2N-1}(n) + g_{NN} \cdot S_1(n) \quad (4n)$$

Such a system of equations can be re-written in a matrix form as:

$$S(n+1) = G \cdot S(n), \quad (5)$$

35

where:

$S(n+1)$ is column matrix containing the state values of the state after a shift,

G is a coefficient matrix comprising the g values indicated in equations 4a-4n, and

$S(n)$ is a column vector of present states.

- 5 Once a state after a shift has been determined, the next state can be calculated using equation (5):

$$S(n+2)=G*S(n+1). \quad (6)$$

- 10 Substituting equation (5) into equation (10) then results into an equation:

$$S(n+2) = G*G*S(n) = G^2*S(n). \quad (7)$$

Further generalization of equation (11) yields an equation:

15
$$S(n+k) = G^k*S(n), \quad (8)$$

where k is a number expressing a state, in which an output is to be computed.

- 20 Applying these principles to FIG. 1, it is obvious that a value of a certain register in next state $S_i(n+1)$ is a function of a value of the preceding register in current state $S_{i-1}(n)$, and -- if a feedback exists -- a value of the output register in current state $S_N(n)$. Consequently, the system of equations (4) will have at most two non-zero coefficients in each of the equations (4a) through (4n).

- 25 As an example, the G matrix for a PN_i shift-register in accordance with FIG. 1b will be developed as follows:

Observing, that there is a connection between stages S_{15} and S_{14} and no feedback from stage S_{15} , it follows that the next state value of S_{15} is equal to previous state value of S_{14} . Thus, equation (4a) will take a form:

30
$$S_{15}(n+1) = 0 \cdot S_{15}(n) + 1 \cdot S_{14}(n) \quad (9)$$

Consequently, the first row of matrix G will contain a non-zero element only in a position g_{12} :

35
$$G_1 = [0100000000000000] \quad (10)$$

Equivalent relation will hold for all stages an input of which is an output of another stage.

Turning to the next stage S_{14} , one can observe that its next state value is equal to previous state value of stage S_{13} summed with a previous state value of stage S_{15} . Thus, the equation (4b) will take a form:

5

$$S_{14}(n+1) = 1 \cdot S_{15}(n) + 1 \cdot S_{13}(n) \quad (11)$$

Consequently, the second row of matrix G will contain a non-zero (unity) element in a position g_{21} and g_{23} :

10

$$G_2 = [1010000000000000] \quad (12)$$

Equivalent relation will hold between all stages an input of which is a sum of outputs of two stages.

15

Reference back to FIG. 3 will expand on these concepts. State memory 212 is initialized to an initial set of states $S_1(n)$, $S_2(n)$, ..., $S_N(n)$. These states are then provided to an output generator 214, and a next state generator 216. Next state generator 216 contains a coefficient matrix G_{ns} formed in accordance with the principles outlined in description of equations (4) and (5). In the exemplary embodiment, the generator polynomial has relatively few feedback taps and, consequently, the resultant matrix G is sparse. This sparseness permits a relatively simple implementation of the matrix operation to be performed using fixed Boolean operator programmed into a field programmable gate array or designed into an application specific integrated circuit (ASIC).

20

Next state generator 216 accepts the set of states $S_1(n)$, $S_2(n)$, ..., $S_N(n)$ from memory 212 to compute a set of new states $S_1(n+K)$, $S_2(n+K)$, ..., $S_N(n+K)$ in accordance with equation (12), and provides the set of new states back to the state memory 212.

30

The output generator 214 performs a matrix operation on the current states in accordance with a matrix G_{os} formed as follows. As explained in description to FIG. 1a, the output of a shift-register is the state $S_N(n)$. From equation (8) follows that:

35

$$S(n+0) = G^0 S(n), \quad (13)$$

where G^0 is a matrix having non-zero elements only in the main diagonal. Inspecting the system of equations (4), it is obvious that value $S_N(n)$ can be calculated using equation (4a). This equation is equivalent to forming a row

matrix G_R by taking the first row of a matrix G_{NS}^0 and multiplying it by a column matrix of states S formed from values $S_1(n), S_2(n), \dots, S_N(n)$. Therefore, the first row of a matrix G_{NS} becomes the last row of matrix G_{OS} . Similarly, from equation (8), the value $S_N(n+1)$ can be calculated by forming a row matrix G_R by taking the first row of a matrix G_{NS}^2 , and multiplying it by a column matrix of states S . Thus, the last row of a matrix G_{NS} becomes the last but one row of matrix G_{OS} . This process of forming the matrix G_{OS} continues until all K rows are filled. In mathematical terms:

$$G_{OS} = \begin{bmatrix} G_{NSL}^K \\ \vdots \\ G_{NSL}^1 \\ G_{NSL}^0 \end{bmatrix}, \quad (14)$$

where G_{NSL}^k is last row of matrix G_{NS}^k .

Once matrix G_{OS} has been formed, the output generator 214 computes the values $S_N(n+1), S_N(n+2), \dots, S_N(n+K)$ by multiplying the matrix G_{OS} by a column matrix of states S :

$$S_N(n+K) = G_{OS} \cdot S(n) \quad (15)$$

A long-code output generator 214 differs from the structure of short-code output generator. The reason is that the long-code generator contains a mask, which can be different for each long-code generator, see, "The cdma2000 ITU-R RTT Candidate Submission" and FIG. 1d. The PN output bit of the long code is a modulo-2 addition of values of the shift registers multiplied by the mask. The output bit can be expressed in matrix notation as follows:

$$pn_{out}(n) = M * S(n), \quad (16)$$

where:

$pn_{out}(n)$ is an output bit in a state n , and
 M is a column mask matrix.

Substituting equation (8) into equation (16) results in:

$$pn_{out}(n+k) = M * G^k * S(n) \quad (17)$$

From equation (10) follows that desired output of $K+1$ parallel bits can be achieved by forming matrix G_{OSL}

$$5 \quad G_{OSL} = \begin{bmatrix} M * G_{NSL}^K \\ \vdots \\ M * G_{NSL}^1 \\ M * G_{NSL}^0 \end{bmatrix}, \quad (18)$$

and, once matrix G_{OSL} has been formed, the output generator 214 computes the values $pn(n)$, $pn(n+1)$, \dots , $pn(n+K)$ by multiplying the matrix G_{OSL} by a column matrix of states S :

$$10 \quad pn(n+K) = G_{OSL} \cdot S(n) \quad (19)$$

At this point of the process the set of states, $S_1(n+K)$, $S_2(n+K)$, \dots , $S_N(n+K)$ is provided to an output generator 214, a next state generator 216, and the whole cycle is repeated.

In particular, let us consider the G matrix for a PN_1 shift-register to be the basic next state generator matrix G_{NS1} :

$$20 \quad G_{NS1} = \begin{bmatrix} 0100000000000000 \\ 1010000000000000 \\ 0001000000000000 \\ 0000100000000000 \\ 0000010000000000 \\ 1000001000000000 \\ 1000000100000000 \\ 1000000010000000 \\ 0000000001000000 \\ 1000000000100000 \\ 0000000000100000 \\ 0000000000010000 \\ 0000000000001000 \\ 0000000000000100 \\ 0000000000000010 \\ 1000000000000000 \end{bmatrix}$$

Matrix G_{NSI}^0 is as follows:

$$G_{NSI0} = \begin{bmatrix} 100000000000000 \\ 010000000000000 \\ 001000000000000 \\ 000100000000000 \\ 000010000000000 \\ 000001000000000 \\ 000000100000000 \\ 000000010000000 \\ 000000001000000 \\ 000000000100000 \\ 000000000010000 \\ 000000000001000 \\ 000000000000100 \\ 000000000000010 \\ 000000000000001 \end{bmatrix}$$

5

Taking the first row of matrix G_{NSI}^0 and last row of matrix G_{NSI} , the matrix G_{OSI2} is formed as follows:

$$G_{OSI2} = \begin{bmatrix} 010000000000000 \\ 100000000000000 \end{bmatrix}$$

10

One ordinarily skilled in the art will recognize that matrix G_{OS} can be modified according to desired PN generator output, without departing from the scope of the invention. For example, if a parallel output $S_N(n)$, $S_N(n+2)$, $S_N(n+4)$, and $S_N(n+6)$ is desired, matrix G_{OS} will comprise in accordance with equation (14) first row of G_{NS}^6 in row one, first row of G_{NS}^4 in row two, first row of G_{NS}^2 in row three, and first row of G_{NS}^0 in row four.

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FIG. 4 depicts a block diagram of a preferred embodiment of the parallel PN generator. In addition to the state memory 212, the output generator 214, and a next state generator 216, it contains a jump generator 218 and a control processor 220. The function of the jump generator 218 is to advance the state by predetermined number of shifts. Such a function is desirable e.g., for forward link acquisition as described in aforementioned U.S. Patent Nos. 5,644,591 and 5,805,648. In the exemplary embodiment, the PN generator is employed in a

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receiver in accordance to an IS-95 standard. The systems designed in accordance with an IS-95 standard comprise base stations utilizing a common PN generator, with a phase offset in increments of 64 chips for a particular pilot signal. Consequently, the jump generator 218 is functionally equivalent to next state generator 216 in that it comprises a coefficient matrix G_{js} formed in accordance with the principles outlined in description of FIG. 1a, and raised to the power of 64.

Next state generator 216 receives the set of states $S_1(n), S_2(n), \dots, S_N(n)$ from memory 212 and generates a set of new states $S_1(n+64), S_2(n+64), \dots, S_N(n+64)$ in accordance with equation (8), and provides the set of new states back to memory 212. The reason for having a separate next state generator 216 and a jump generator 218 is that in general $K \neq L$, and, consequently, the matrices G_{os} and G_{js} are different. As described above, the present invention is preferably implemented in hardware adapted to the specific operation and designed to perform a specific task.

The function of the control processor 220 is to coordinate cooperation between the different subsystems, and to control bit stuffing. As described, the short-code PN sequences have a period of 2^{15} generating polynomials, and from them derived matrices, generate only sequences with period $2^{15} - 1$. The control processor 200 monitors the output of the next state generator 216 for the state preceding the state corresponding to a period $2^{15} - 1$, for which a computation of next state according to equation (8) would exceed the state corresponding to a period $2^{15} - 1$. Once the control processor 200 detects such state it performs two operations. It will cause the output generator 214 to compute the output state values, and overwrites the last output state value with '0'. It will then avoid writing the output of the next state generator 216 into state memory 212, and will initialize the state memory 212 to initial set of states $S_1(n), S_2(n), \dots, S_N(n)$.

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FIG. 5 depicts a simplified block diagram of an exemplary receiver chain using PN generators in accordance with the invention. The RF signal arriving at the antenna 400 is provided to the receiver (RCVR) 402, which downconverts the received signal to a baseband frequency, producing I and Q components of the signal. These components are simultaneously provided to a searcher 404 and demodulators 406a, \dots , 406c. The task of the searcher 404 is to perform searches in code space to identify candidate signals to be added to the Active Set of the remote station in order to maximize the quality of the received signal. To accomplish this task, searcher 404 will control parameters of the PN

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sequences generators, devised in accordance with the principles outlined in present invention. An exemplary method for performing acquisition and searching in a CDMA communication system is described in detail in aforementioned U.S. Patent Nos. 5,644,591 and 5,805,648

5 In order to be effective, a receiver must be able to operate in a multipath environment and must be able to adapt to changes in physical location. In the aforementioned U.S. Patent Nos. 5,101,501 and 5,109,390, a method for exploiting the reception of multiple version of a signal is described. Demodulators 406a, 406b and 406c demodulate redundant versions of the same
10 signal. These redundant version either correspond to multipath propagations of a signal from a single source or from multiple transmissions of the same information from multiple base stations in a soft handoff condition.

The demodulated signals from demodulators 406a, . . . , 406c are provided to combiner 410, which combines the signals and provides them for
15 further processing to a de-interleaver 412 and decoder 414.

FIG. 6 illustrates the exemplary embodiment of the receiver structure of the present invention. The signal is received at antenna 400 and provided to receiver (RCVR) 402. Receiver 402 down converts, amplifies, filters, and
20 samples the received signal, and provides digital samples to buffer 402. In response to signals from control processor 403, a selected set of samples from buffer 404 are provided to despreader 408. In addition, in response to a signal from control processor 403, PN generator 406 provides a portion of a PN sequence to despreader 408.

25 Despreader 408 despreads the signal in accordance with the portion of the PN sequence provided by PN generator 406 which operates in accordance with the present invention. Within despreader 408 the PN sequence is provided to pilot despreader 412, which despreads the received signal in accordance with the portion of the short PN sequence provided by PN
30 generator 406 and the Walsh covering sequence for the pilot signal. In the exemplary embodiment, the pilot signal is covered with the Walsh zero sequence and as such does not effect the despreading operation performed by pilot despreader 412. In addition, the portion of the short PN sequence is provided to traffic despreader 414, which despreads the signal in accordance
35 with the short PN sequence and the Walsh traffic covering sequence W_T .

The result of the despreading operation performed by pilot despreader 412 and the result of the despreading operation performed by traffic despreader 414 are provided to dot product circuit 414. The pilot signal has known symbols and can be used to remove the phase ambiguities introduced by the

propagation path as described in the aforementioned U.S. Patent No. 5,506,865. The result of the dot product operation is provided to combiner 410. Combiner 410 combines redundantly despread version of the same symbols whether transmitted by different base stations in a soft handoff environment or by the same base station traversing different propagation paths in a multipath environment.

In accordance with an exemplary demodulation chain embodiment, and previous discussion follows that a first set of matrices is required for the short-code PN generator for the I component 516, a second set for the short-code PN generator for the Q component 518, and a third set for the long-code PN generator 504.

1. Acquisition mode.

In the exemplary embodiment, the receiver is able to rapidly determine jump 64 chips ahead in the PN sequence in order to perform a correlation process to determine the correlation energy of between the received signal and a portion of the PN sequence.

In the generation of the short PN₁ sequence, state memory 212 provides the current state of the PN sequence S(n) to next state generator 216. Next state generator 216 generates the state of the PN sequence S(n+2) two cycles in advance by left-multiplying the PN sequence S(n) by the matrix G_{NSI2}:

$$G_{NSI2} = \begin{bmatrix} 10100000000000 \\ 01010000000000 \\ 00001000000000 \\ 00000100000000 \\ 10000010000000 \\ 11000001000000 \\ 11000000100000 \\ 01000000010000 \\ 10000000001000 \\ 01000000000100 \\ 00000000000010 \\ 00000000000001 \\ 00000000000000 \\ 10000000000000 \\ 01000000000000 \end{bmatrix}$$

In the generation of the short PN_i sequence, state memory 212 provides the current state of the PN sequence S(n) to jump generator 218. Jump generator 218 generates the state of the PN sequence S(n+2) sixty-four (64) cycles in advance by left-multiplying the PN sequence S(n) by the matrix G_{JS164} :

$$G_{JS164} = \begin{bmatrix} 101011010100101 \\ 010101101010010 \\ 000001100001100 \\ 000000110000110 \\ 000000011000011 \\ 000000001100001 \\ 101011010010101 \\ 011110111101111 \\ 000100001010010 \\ 000010000101001 \\ 101010010110001 \\ 110101001011000 \\ 011010100101100 \\ 101101010010110 \\ 010110101001011 \end{bmatrix}$$

In the generation of the short PN_i sequence, the next state generator 216 or the jump generator 218 provides the current state of the PN sequence S(n) to output generator 214. Output generator 214 computes the values S_N(n+1), S_N(n+2), . . . , S_N(n+K) left-multiplying a column matrix of states S(n) by the matrix G_{OS12} :

$$G_{OS12} = \begin{bmatrix} 0100000000000000 \\ 1000000000000000 \end{bmatrix}$$

The short-code PN generator for the Q component 518 uses an algorithm for PN sequence generation, identical to the one for the acquisition mode. Consequently, the set of matrices as well as their application is identical.

$$G_{NSQ2} = \begin{bmatrix} 001000000000000 \\ 100100000000000 \\ 110010000000000 \\ 110001000000000 \\ 010000100000000 \\ 000000010000000 \\ 000000001000000 \\ 100000000100000 \\ 110000000010000 \\ 110000000001000 \\ 110000000000100 \\ 010000000000010 \\ 000000000000001 \\ 100000000000000 \\ 010000000000000 \end{bmatrix}$$

$$G_{JSQ64} = \begin{bmatrix} 100011001011100 \\ 010001100101110 \\ 001000110010111 \\ 000111010010111 \\ 100000100010111 \\ 110011011010111 \\ 011001101101011 \\ 101100110110101 \\ 110110011011010 \\ 011000000110001 \\ 101111001000100 \\ 110100101111110 \\ 011001011100011 \\ 001100101110001 \\ 000110010111000 \end{bmatrix}$$

5

$$G_{osq2} = \begin{bmatrix} 010000000000000 \\ 100000000000000 \end{bmatrix}$$

10 In the generation of the long-code PN sequence, state memory 212 provides the current state of the PN sequence $S(n)$ to next state generator 216.

Next state generator 216 generates the state of the PN sequence $S(n+2)$ two cycles in advance by left-multiplying the PN sequence $S(n)$ by the matrix G_{NSL2} :

5 In the generation of the long-code PN sequence, state memory 212 provides the current state of the PN sequence $S(n)$ to jump generator 218. Jump generator 218 generates the state of the PN sequence $S(n+64)$ sixty-four (64) cycles in advance by left-multiplying the PN sequence $S(n)$ by the matrix G_{JSL64} :

$G_{JSL64} =$

```

01110010001101110111101111100100001111110
10111001000110111011110111111001000011111
01011100100011011101111011111001000011111
00101110010001101110111101111100100001111
10010111001000110111011110111110010000111
11001011100100011011101111011111001000011
01100101110010001101110111101111100100001
01000000110100110001010100000101111101110
00100000011010011000101010000010111110111
011000100000001110111110101100110110000101
001100010000000111011111010110011011000010
111010101011011110010100010111101100011111
011101010101101111001010001011110110001111
101110101010110111100101000101111011000111
010111010101011011110010100010111101100011
01011100100111000000001010110111111001111
010111000111100101111010101010011110011001
010111000000101111000110101001101110110010
001011100000010111100011010100110111011001
100101110000001011110001101010011011101100
001110011011011000000011001001101100001000
011011101110110001111010011000010111111010
101101110111011000111101001100001011111101
001010011000110001100101011010100100000000
111001101111000101001001010001110011111110
000000010100111111011111010100011000000001
011100101001000010010100010110101101111110
001110010100100001001010001011010110111111
100111001010010000100101000101101011011111
110011100101001000010010100010110101101111
011001110010100100001001010001011010110111
001100111001010010000100101000101101011011
011010111111110100111001101000110111010011
101101011111111010011100110100011011101001
010110101111111101001110011010001101110100
110111110100100011011100110001100111000100
100111011001001100010101100100010010011100
0011110011111111011110001001110101000110000
1001111001111111101111000100111010100011000
101111010000100011000111101111001011110010
001011001011001100011000001011000100000111
11100100011011101111011111001000011111101

```

In the generation of the long-code PN sequence, the next state generator 216 or the jump generator 218 provides the current state of the PN sequence $S(n)$ to output generator 214. Output generator 214 first computes the output
5 state matrix G_{OSL} by left-multiplying matrix M by matrices G_{NSLO} :

,and by matrix G_{NS1} :

, and then computes the output bits $pn_{out}(n+k)$ by multiplying the resulting matrix G_{osl} by a column matrix of states S .

2. Demodulation mode:

5 The demodulation mode uses algorithm for PN sequence generation, identical to the one for the acquisition mode. Consequently, the set of matrices as well as their application is identical.

The short-code PN generator for the I component 516 comprises the following matrices:

10

$$G_{NSI8} = \begin{bmatrix} 010010101000000 \\ 001001010100000 \\ 110110000010000 \\ 111011000001000 \\ 011101100000100 \\ 101110110000010 \\ 000101110000001 \\ 010000010000000 \\ 011010100000000 \\ 001101010000000 \\ 010100000000000 \\ 101010000000000 \\ 010101000000000 \\ 001010100000000 \\ 100101010000000 \end{bmatrix}$$

$$G_{JSI64} = \begin{bmatrix} 101011010100101 \\ 010101101010010 \\ 000001100001100 \\ 000000110000110 \\ 000000011000011 \\ 000000001100001 \\ 101011010010101 \\ 011110111101111 \\ 000100001010010 \\ 000010000101001 \\ 101010010110001 \\ 110101001011000 \\ 011010100101100 \\ 101101010010110 \\ 010110101001011 \end{bmatrix}$$

5

$$G_{Q518} = \begin{bmatrix} 1001010100000000 \\ 0010101000000000 \\ 0101010000000000 \\ 1010100000000000 \\ 0101000000000000 \\ 1010000000000000 \\ 0100000000000000 \\ 1000000000000000 \end{bmatrix}$$

The short-code PN generator for the Q component 518 comprises the following matrices:

$$G_{NSQ8} = \begin{bmatrix} 101111001000000 \\ 010111100100000 \\ 101011110010000 \\ 011010110001000 \\ 000010010000100 \\ 001110000000010 \\ 100111000000001 \\ 110011100000000 \\ 111001110000000 \\ 010011110000000 \\ 000110110000000 \\ 101100010000000 \\ 111001000000000 \\ 111100100000000 \\ 011110010000000 \end{bmatrix}$$

$$G_{JSQ64} = \begin{bmatrix} 100011001011100 \\ 010001100101110 \\ 001000110010111 \\ 000111010010111 \\ 100000100010111 \\ 110011011010111 \\ 011001101101011 \\ 101100110110101 \\ 110110011011010 \\ 011000000110001 \\ 101111001000100 \\ 110100101111110 \\ 011001011100011 \\ 001100101110001 \\ 000110010111000 \end{bmatrix}$$

$$G_{os08} = \begin{bmatrix} 0111100100000000 \\ 1111001000000000 \\ 1110010000000000 \\ 1100100000000000 \\ 1001000000000000 \\ 0010000000000000 \\ 0100000000000000 \\ 1000000000000000 \end{bmatrix}$$

The long-code PN generator for 518 comprises the following matrices:

$G_{JSL64} =$

```
011100100011011101111011111100100001111110
101110010001101110111101111110010000111111
010111001000110111011110111111001000011111
001011100100011011101111011111100100001111
100101110010001101110111101111110010000111
110010111001000110111011110111111001000011
011001011100100011011101111011111100100001
01000000110100110001010100000101111101110
00100000011010011000101010000010111110111
011000100000001110111110101100110110000101
001100010000000111011111010110011011000010
111010101011011110010100010111101100011111
011101010101101111001010001011110110001111
101110101010110111100101000101111011000111
010111010101011011110010100010111101100011
010111001001110000000010101101111111001111
010111000111100101111010101010011110011001
010111000000101111000110101001101110110010
001011100000010111100011010100110111011001
100101110000001011110001101010011011101100
001110011011011000000011001001101100001000
01101110111011000111101001100001011111010
10110111011101100011110100110000101111101
001010011000110001100101011010100100000000
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000000010100111111011111010100011000000001
01110010100100001001010001011010110111110
00111001010010000100101000101101011011111
10011100101001000010010100010110101101111
11001110010100100001001010001011010110111
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001100111001010010000100101000101101011011
011010111111110100111001101000110111010011
101101011111111010011100110100011011101001
010110101111111101001110011010001101110100
110111110100100011011100110001100111000100
100111011001001100010101100100010010011100
001111001111111101111000100111010100011000
100111100111111101111000100111010100011000
101111010000100011000111101111001011110010
001011001011001100011000001011000100000111
111001000110111011111011111100100001111101
```


$$G_{OSL88} =$$
[illegible]

The previous description of the preferred embodiments is provided to enable any person skilled in the art to make or use the present invention. The various modifications to these embodiments will be readily apparent to those skilled in the art, and the generic principles defined herein may be applied to other embodiments without the use of the inventive faculty. Thus, the present invention is not intended to be limited to the embodiments shown herein but is to be accorded the widest scope consistent with the principles and novel features disclosed herein.

WHAT IS CLAIMED IS:

1. An apparatus for generating multiple bits of a pseudonoise sequence with each clock pulse by computing the bits in parallel, comprising:
 - a) a state memory;
 - b) a next state generator communicatively connected with said state memory; and
 - c) an output generator communicatively connected with said state memory and said next state generator.
2. The apparatus of claim 1 wherein said state memory has been configured to hold:
 - a) a set of initial values of states; and
 - b) a set of values of states generated by said next state generator or a jump generator.
3. The apparatus of claim 1 wherein said set of initial values of states comprises:
 - a) coefficients of a generating polynomial.
4. The apparatus of claim 3 wherein said generating polynomial is:

$$P_1(x) = x^{15} + x^{13} + x^9 + x^8 + x^7 + x^5 + 1$$
5. The apparatus of claim 3 wherein said generating polynomial is:

$$P_0(x) = x^{15} + x^{12} + x^{11} + x^{10} + x^6 + x^5 + x^4 + x^3 + 1$$
6. The apparatus of claim 3 wherein said generating polynomial is:

$$P(x) = x^{42} + x^{35} + x^{33} + x^{31} + x^{27} + x^{26} + x^{25} + x^{22} + x^{21} + x^{19} + x^{18} + x^{17} + x^{16} + (1) + x^{10} + x^7 + x^6 + x^5 + x^3 + x^2 + x + 1.$$

2 7. The apparatus of claim 1 wherein said next state generator has been configured to:

- 4 a) accept one set of values of states;
 6 b) generate another set of values of states a first pre-determined number of clocks apart from current state by multiplying said accepted values by a next step matrix; and
 8 c) provide said another set of values of states to said memory and said output generator.

2 8. The apparatus of claim 7 wherein said first pre-determined number of clocks is two and said next step matrix G_{NS12} is:

$$G_{NS12} = \begin{bmatrix} 10100000000000 \\ 01010000000000 \\ 00001000000000 \\ 00000100000000 \\ 10000010000000 \\ 11000001000000 \\ 11000000100000 \\ 01000000010000 \\ 10000000001000 \\ 01000000000100 \\ 00000000000010 \\ 00000000000001 \\ 00000000000000 \\ 10000000000000 \\ 01000000000000 \end{bmatrix}$$

4

2 9. The apparatus of claim 7 wherein said first pre-determined number of clocks is two and said next step matrix G_{NSQ2} is:

$$G_{NSQ2} = \begin{bmatrix} 00100000000000 \\ 10010000000000 \\ 11001000000000 \\ 11000100000000 \\ 01000010000000 \\ 00000001000000 \\ 00000000100000 \\ 10000000010000 \\ 11000000001000 \\ 11000000000100 \\ 11000000000010 \\ 01000000000001 \\ 00000000000001 \\ 10000000000000 \\ 01000000000000 \end{bmatrix}$$

4

- 2 10. The apparatus of claim 7 wherein said a first pre-determined number of
clocks is eight and said next step matrix G_{NS18} is:

4

$$G_{NS18} = \begin{bmatrix} 01001010100000 \\ 00100101010000 \\ 11011000001000 \\ 11101100000100 \\ 01110110000010 \\ 10111011000001 \\ 00010111000001 \\ 01000001000000 \\ 01101010000000 \\ 00110101000000 \\ 01010000000000 \\ 10101000000000 \\ 01010100000000 \\ 00101010000000 \\ 10010101000000 \end{bmatrix}$$

- 2 11. The apparatus of claim 7 wherein said a first pre-determined number of
clocks is eight and said next step matrix G_{NSQ2} is:

$$G_{NSQ8} = \begin{bmatrix} 101111001000000 \\ 010111100100000 \\ 101011110010000 \\ 011010110001000 \\ 000010010000100 \\ 001110000000010 \\ 100111000000001 \\ 110011100000000 \\ 111001110000000 \\ 010011110000000 \\ 000110110000000 \\ 101100010000000 \\ 111001000000000 \\ 111100100000000 \\ 011110010000000 \end{bmatrix}$$

4

2 12. The apparatus of claim 1 wherein said output generator has been
configured to:

- 4 a) one set of values of states; and
b) generate multiple output bits in parallel by multiplying said
6 accepted values by an output state matrix.

2 13. The apparatus of claim 12 wherein said multiple is two and said output
state matrix G_{OS12} is:

4

$$G_{OS12} = \begin{bmatrix} 010000000000000 \\ 100000000000000 \end{bmatrix}$$

2 14. The apparatus of claim 12 wherein said multiple is two and said output
state matrix G_{OSQ2} is:

4

$$G_{OSQ2} = \begin{bmatrix} 010000000000000 \\ 100000000000000 \end{bmatrix}$$

- 2 15. The apparatus of claim 12 wherein said multiple is eight and said output
state matrix G_{OS18} is:

4

$$G_{OS18} = \begin{bmatrix} 1001010100000000 \\ 0010101000000000 \\ 0101010000000000 \\ 1010100000000000 \\ 0101000000000000 \\ 1010000000000000 \\ 0100000000000000 \\ 1000000000000000 \end{bmatrix}$$

- 2 16. The apparatus of claim 12 wherein said multiple is eight and said output
state matrix G_{OSQ8} is:

4

$$G_{OSQ8} = \begin{bmatrix} 0111100100000000 \\ 1111001000000000 \\ 1110010000000000 \\ 1100100000000000 \\ 1001000000000000 \\ 0010000000000000 \\ 0100000000000000 \\ 1000000000000000 \end{bmatrix}$$

- 2 17. The apparatus of claim 1 further comprising a jump generator.

- 2 18. The apparatus of claim 17 wherein said jump generator has been
configured to:

- 4 a) accept one set of values of states;
b) generate values of states a second pre-determined number of
6 clocks apart from current state by multiplying said accepted values by a jump
state matrix; and
8 c) provide said values of states to said memory and said output
generator.

- 2 19. The apparatus of claim 18 wherein said second pre-determined number is sixty-four and said jump state matrix G_{JS164} is:

$$G_{JS164} = \begin{bmatrix} 101011010100101 \\ 010101101010010 \\ 000001100001100 \\ 000000110000110 \\ 000000011000011 \\ 000000001100001 \\ 101011010010101 \\ 011110111101111 \\ 000100001010010 \\ 000010000101001 \\ 101010010110001 \\ 110101001011000 \\ 011010100101100 \\ 101101010010110 \\ 010110101001011 \end{bmatrix}$$

4

- 2 20. The apparatus of claim 18 wherein said second pre-determined number is sixty-four and said jump state matrix G_{JSQ64} is:

$$G_{JSQ64} = \begin{bmatrix} 100011001011100 \\ 010001100101110 \\ 001000110010111 \\ 000111010010111 \\ 100000100010111 \\ 110011011010111 \\ 011001101101011 \\ 101100110110101 \\ 110110011011010 \\ 011000000110001 \\ 101111001000100 \\ 110100101111110 \\ 011001011100011 \\ 001100101110001 \\ 000110010111000 \end{bmatrix}$$

21. The apparatus of claim 1 further comprising a controller.

22. The apparatus of claim 21 wherein said controller has been configured to
 2 monitor output bits of said next state generator for a pre-determined
 combination, and when said pre-determined combination has been reached to:
- 4 a) overwrite an appropriate output bit value with a value of '0';
 - 6 b) void writing values of states generated by said next state
 generator to said state memory; and
 - 8 c) instruct said state memory to provide a set of initial values of
 states to said next state generator.
23. A pseudonoise (PN) sequence generator comprising:
- 2 a) state memory for storing at least one state of a PN generator
 polynomial;
 - 4 b) next state generator for receiving said at least one state of said PN
 generator polynomial and for generating a second state of said PN generator
 polynomial by performing a matrix operation upon said at least one state of
 said PN generator polynomial; and
 - 8 c) output generator for receiving said at least one state of said PN
 generator polynomial and for generating at least one PN sequence output by
 10 performing a matrix operation upon said at least one state of said PN generator
 polynomial.
24. The apparatus of Claim 23 wherein said at least one state comprises the
 2 fifteen component state of a PN short code.
25. The apparatus of Claim 23 wherein said at least one state comprises the
 2 forty two component state of a PN long code.

2
$$P_Q(x) = x^{15} + x^{12} + x^{11} + x^{10} + x^6 + x^5 + x^4 + x^3 + 1$$

28. The apparatus of Claim 23 wherein said generator polynomial (P) is:
- 2
$$P(x) = x^{42} + x^{35} + x^{33} + x^{31} + x^{27} + x^{26} + x^{25} + x^{22} + x^{21} + x^{19} + x^{18} + x^{17} + x^{16} +$$

$$(1) + x^{10} + x^7 + x^6 + x^5 + x^3 + x^2 + x + 1.$$

29. The apparatus of Claim 23 wherein said next state generator computes
 2 the state of PN sequence generator two clock cycles in the future and performs
 said matrix operation in accordance with the matrix GNSI21:

$$G_{NSI2} = \begin{bmatrix} 101000000000000 \\ 010100000000000 \\ 000010000000000 \\ 000001000000000 \\ 100000100000000 \\ 110000010000000 \\ 110000001000000 \\ 010000000100000 \\ 100000000010000 \\ 010000000001000 \\ 000000000000100 \\ 000000000000010 \\ 000000000000001 \\ 100000000000000 \\ 010000000000000 \end{bmatrix}.$$

4

30. The apparatus of Claim 23 wherein said next state generator performs
 2 said matrix operation in accordance with the matrix GNSQ2:

$$G_{NSQ2} = \begin{bmatrix} 001000000000000 \\ 100100000000000 \\ 110010000000000 \\ 110001000000000 \\ 010000100000000 \\ 000000010000000 \\ 000000001000000 \\ 100000000100000 \\ 110000000010000 \\ 110000000001000 \\ 110000000000100 \\ 010000000000010 \\ 000000000000001 \\ 100000000000000 \\ 010000000000000 \end{bmatrix}.$$

31. The apparatus of Claim 23 wherein said output generator computes the next two outputs of said PN sequence generator and performs said matrix operation in accordance with the matrix GOSI2:

$$G_{OSI2} = \begin{bmatrix} 0100000000000000 \\ 1000000000000000 \end{bmatrix}.$$

32. The apparatus of Claim 23 wherein said PN generator programmed into an ASIC.

33. The apparatus of Claim 23 wherein said PN generator programmed into a field programmable gate array.

34. A method for generating multiple bits of a pseudonoise sequence with each clock pulse by computing the bits in parallel, comprising the steps of:

- a) storing at least one set of values of states in a state memory;
- b) generating a second set of values of states by a next state generator, said second set being derived from said at least one set; and
- c) generating a set of output bits in parallel by an output generator, said set of output bits being derived from said at least one set of values of states.

35. The method of claim 34, wherein the step of storing at least one set of values of states comprises the steps of:

- a) holding a set of initial values of states; and
- b) holding another set of values of states from said next state generator or from a jump generator.

36. The method of claim 34, wherein the step of generating a second set of values of states comprises the step of:

- a) multiplying said at least one set of values of states by a next step matrix.

37. The method of claim 34, wherein the step of generating a set of output bits in parallel comprises the step of:

- a) multiplying said at least one set of values of states by an output state matrix.

38. The method of claim 34, further comprising the step of monitoring a set
2 of values of states of said next state generator for a pre-determined
combination.

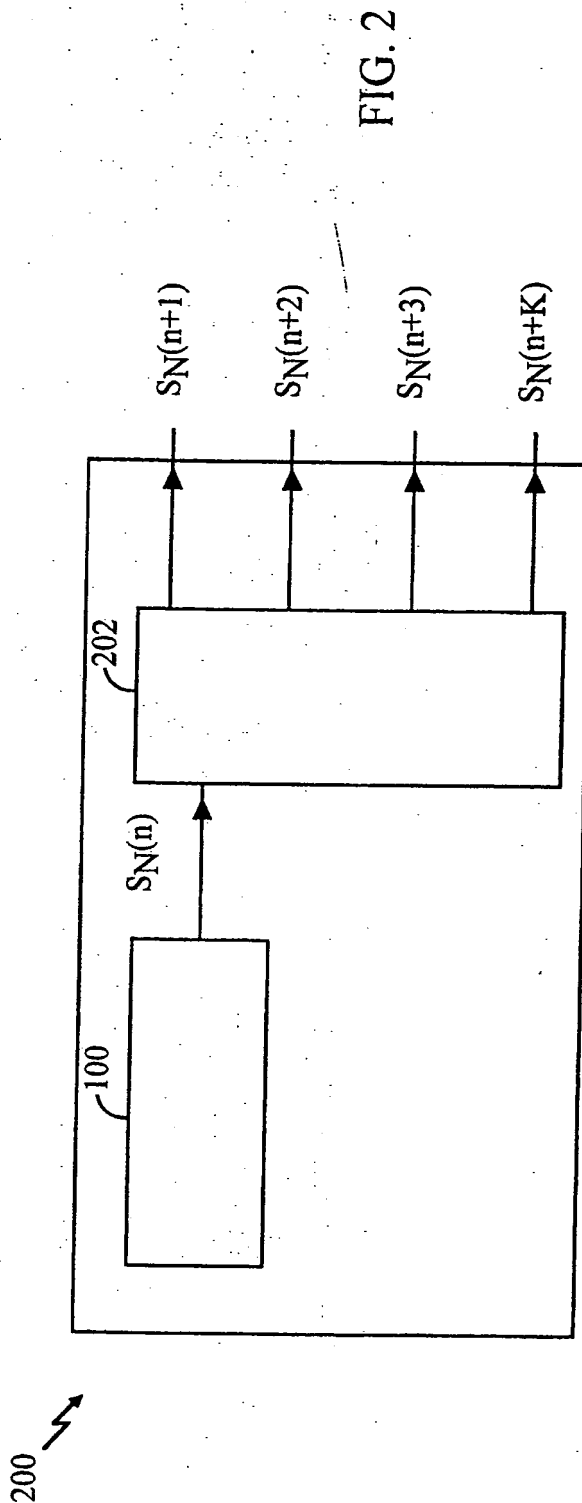
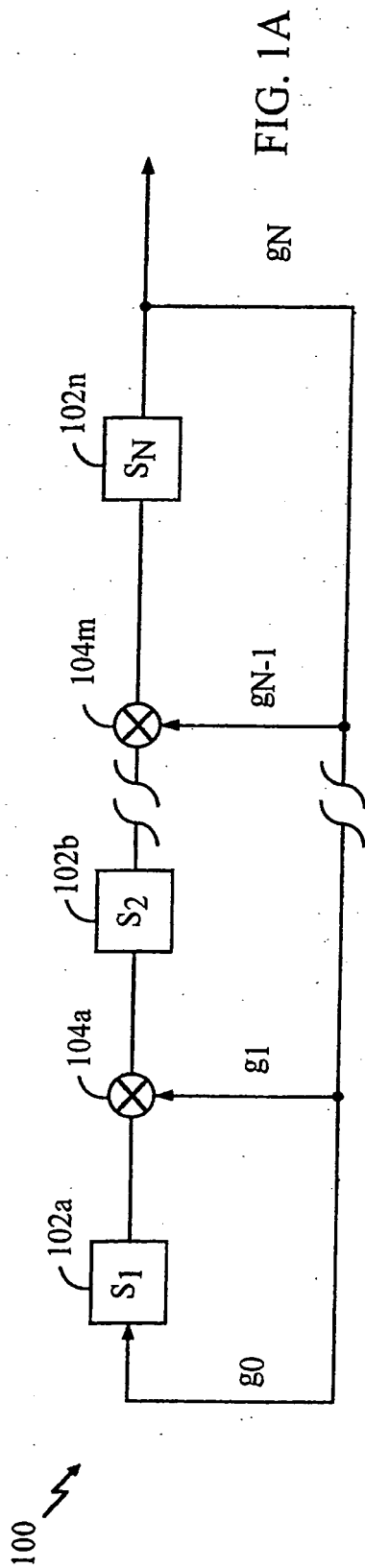
39. The method of claim 38, wherein upon detecting said pre-determined
2 combination, the method further comprises the steps of:

- a) overwriting an appropriate output bit value with a value of '0';
- 4 b) voiding writing said second set of values of states generated by
said next state generator to said state memory; and
- 6 c) instructing said state memory to provide a set of initial values of
states to said next state generator.

40. The method of claim 34, further comprising the step of generating a third
2 set of values of states by a jump state generator, said second set being derived
from said at least one set.

41. The method of claim 40, wherein the step of generating a third set of
2 values of states by a jump state generator comprises the step of:

- a) multiplying said at least one set of values of states by a jump state
4 matrix.



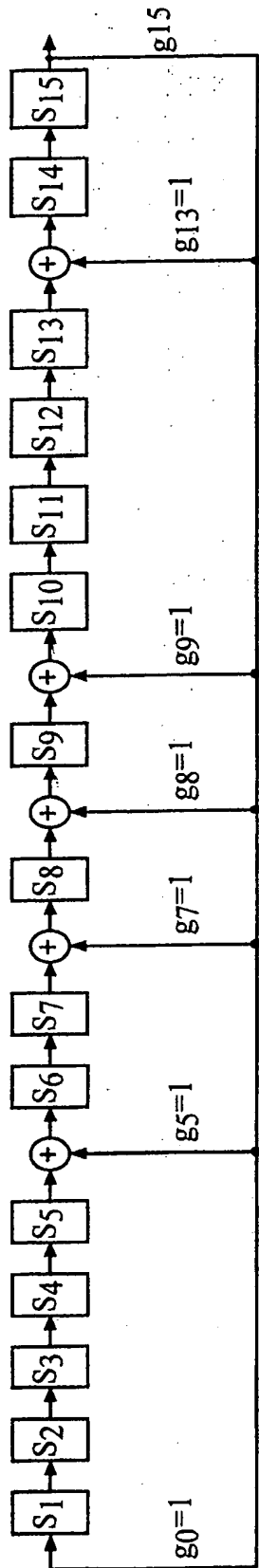


FIG. 1B

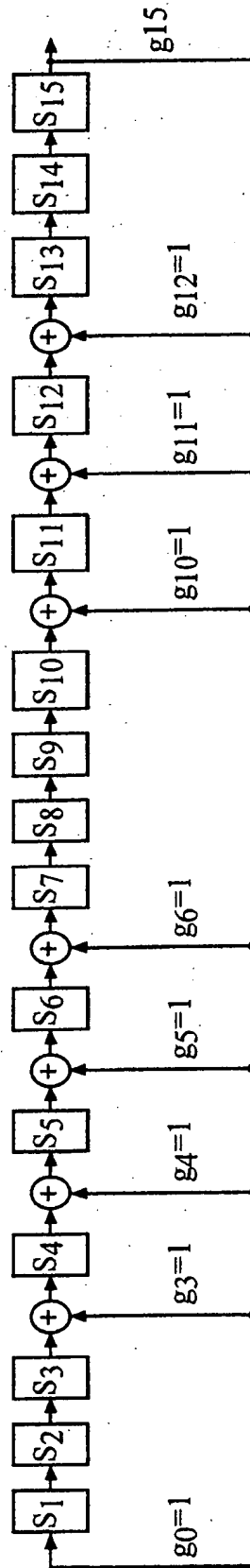


FIG. 1C

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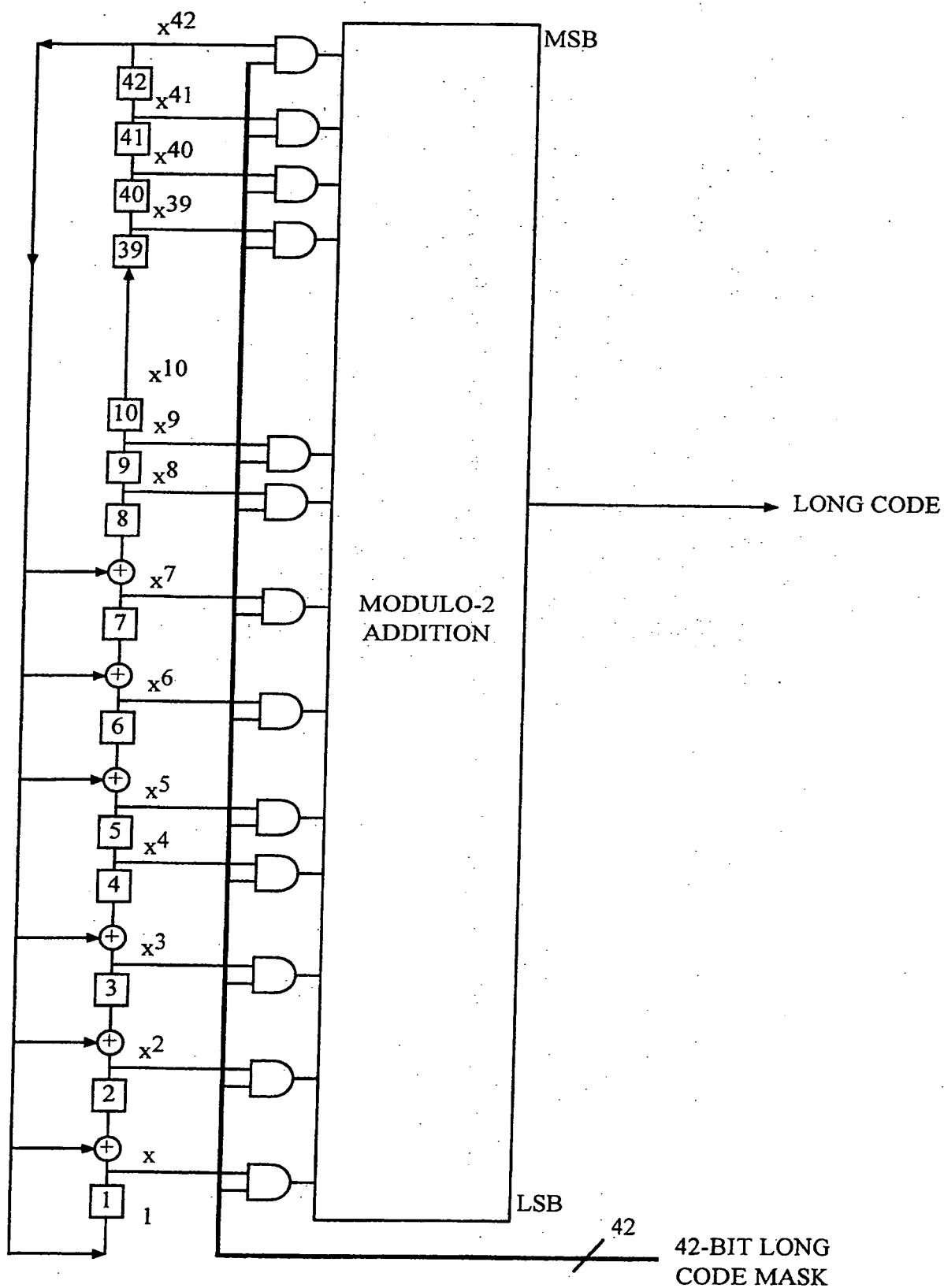


FIG. 1D

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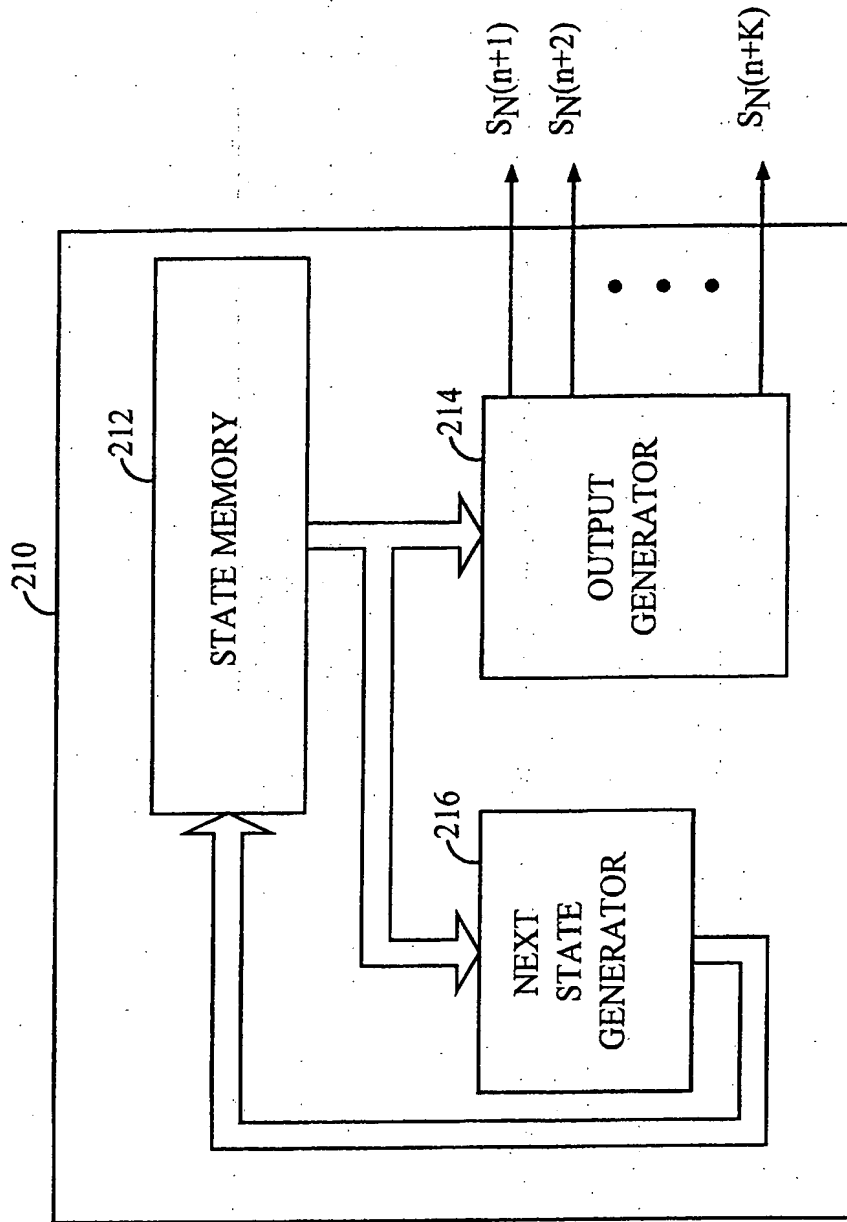


FIG. 3

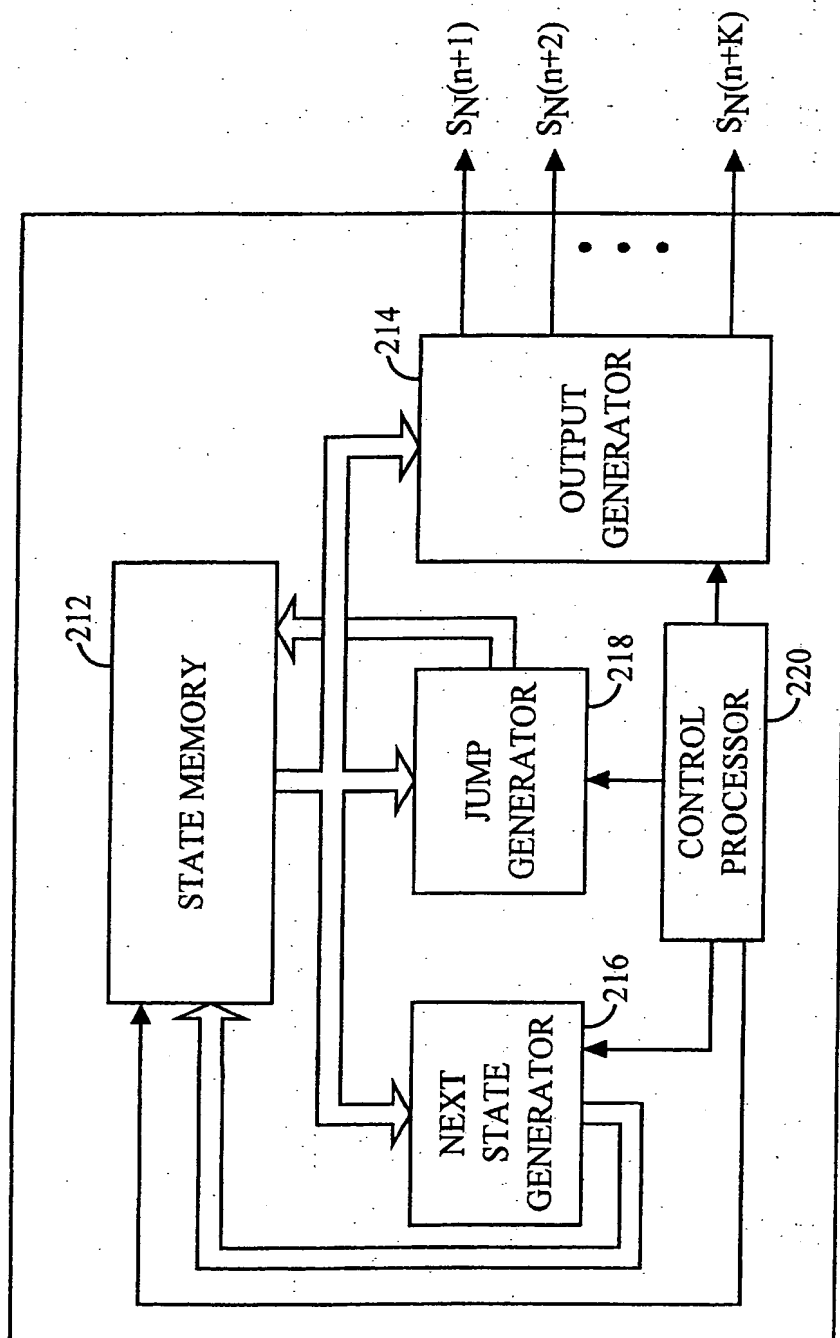


FIG. 4

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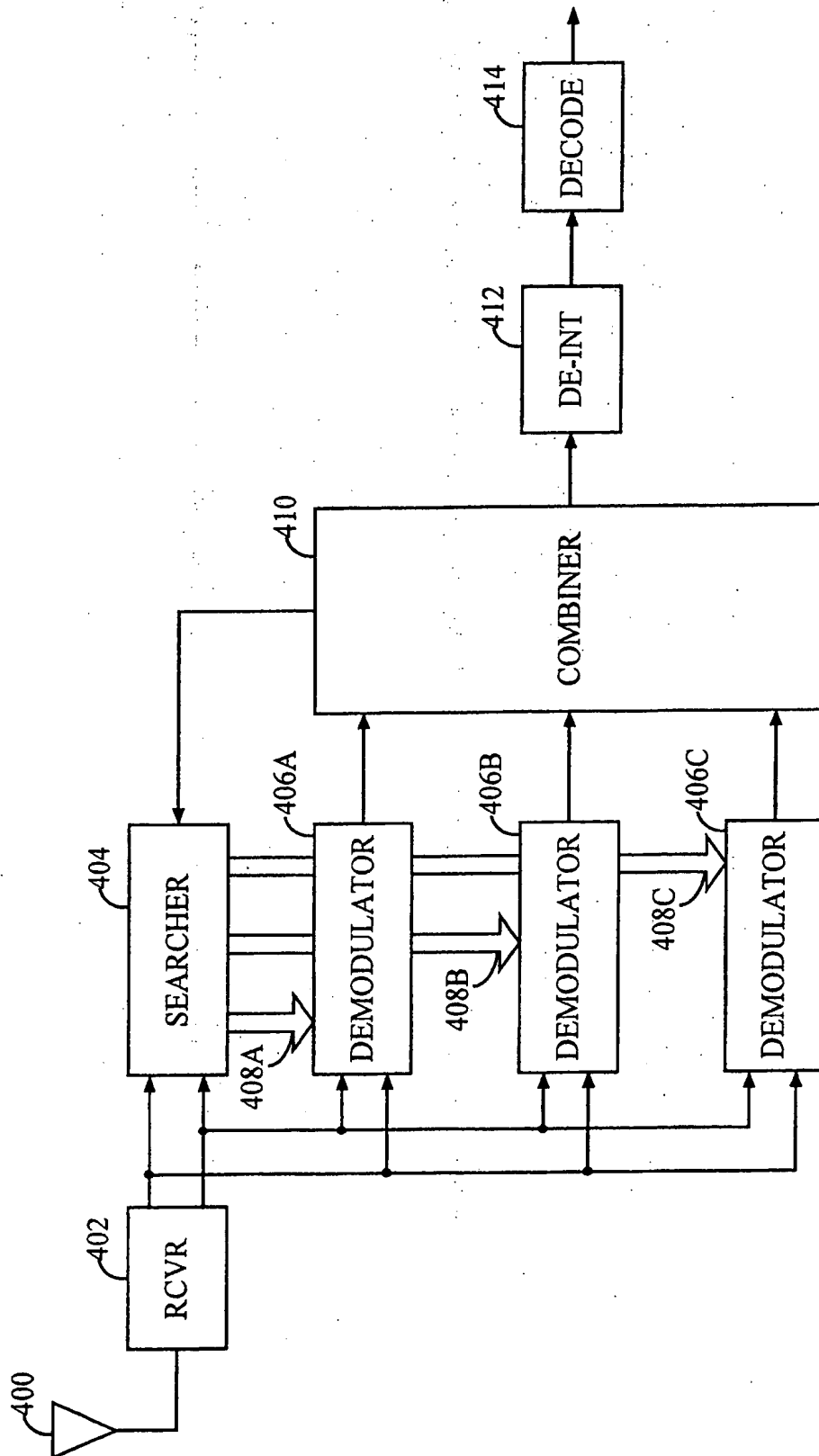


FIG. 5

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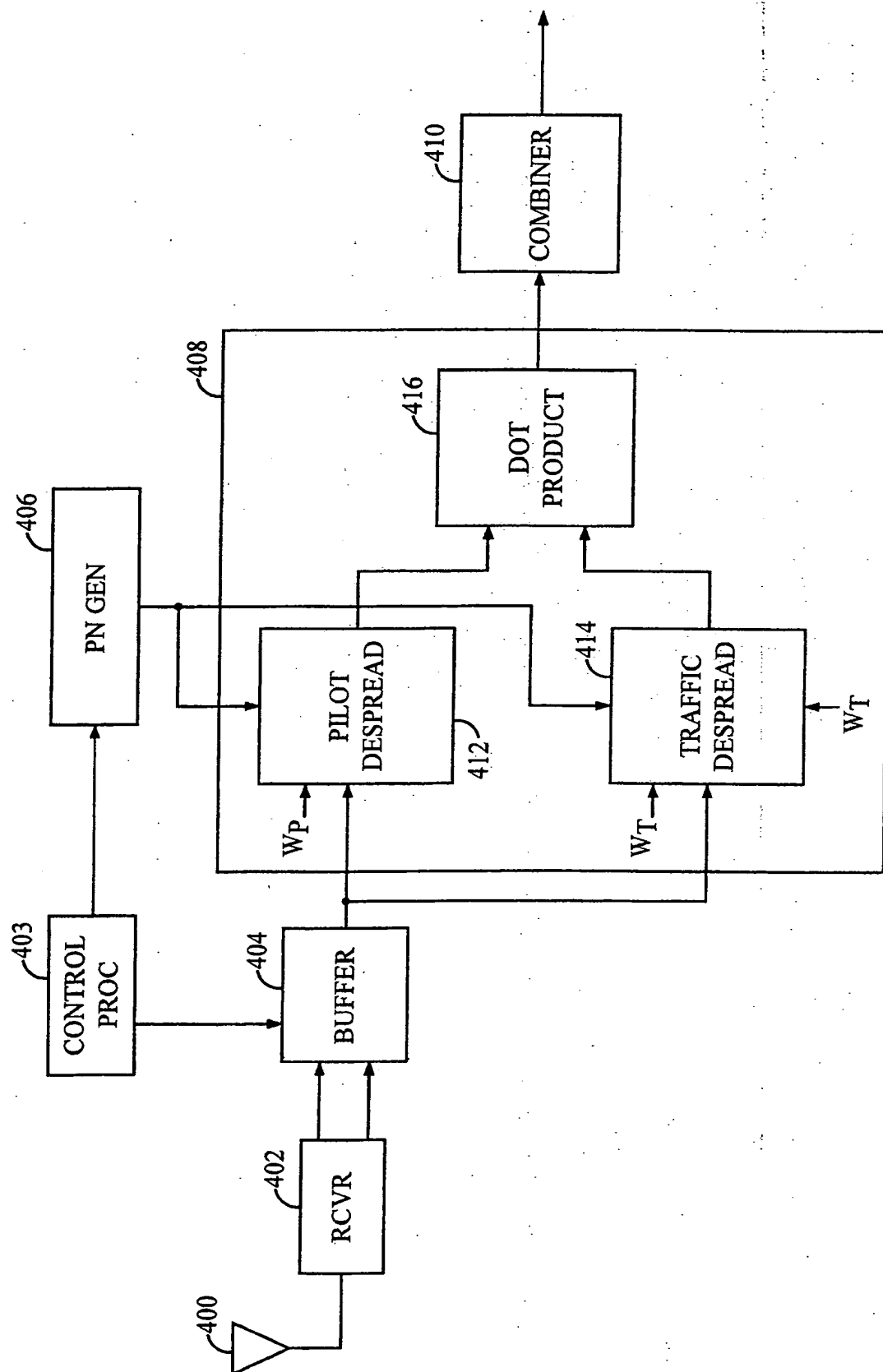


FIG. 6

INTERNATIONAL SEARCH REPORT

Intern. Application No
PCT/US 00/23949

A. CLASSIFICATION OF SUBJECT MATTER

IPC 7 G06F7/58

According to International Patent Classification (IPC) or to both national classification and IPC

B. FIELDS SEARCHED

Minimum documentation searched (classification system followed by classification symbols)
IPC 7 G06F H03K

Documentation searched other than minimum documentation to the extent that such documents are included in the fields searched

Electronic data base consulted during the international search (name of data base and, where practical, search terms used)

EPO-Internal, PAJ, WPI Data

C. DOCUMENTS CONSIDERED TO BE RELEVANT

Category *	Citation of document, with indication, where appropriate, of the relevant passages	Relevant to claim No.
X	O. KAESTNER: "Implementing Branch instructions with Polynomial Counters" COMPUTER DESIGN., vol. 14, no. 1, January 1975 (1975-01), pages 69-75, XP002155211 PENNWELL PUBL. LITTLETON, MASSACHUSETTS., US ISSN: 0010-4566 the whole document	1-41
X	US 3 881 099 A (AILETT CLAUDE ET AL) 29 April 1975 (1975-04-29) figures	1-41
X	US 5 910 907 A (BRADLEY ALAN S ET AL) 8 June 1999 (1999-06-08) figures	1-41
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☒ Further documents are listed in the continuation of box C.

☒ Patent family members are listed in annex.

* Special categories of cited documents:

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- *X* document of particular relevance; the claimed invention cannot be considered novel or cannot be considered to involve an inventive step when the document is taken alone
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Date of the actual completion of the international search

12 December 2000

Date of mailing of the international search report

29/12/2000

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INTERNATIONAL SEARCH REPORT

International Application No
PCT/US 00/23949

C.(Continuation) DOCUMENTS CONSIDERED TO BE RELEVANT		
Category *	Citation of document, with indication, where appropriate, of the relevant passages	Relevant to claim No.
A	US 5 228 054 A (GILHOUSEN KLEIN S ET AL) 13 July 1993 (1993-07-13) column 4, line 37 - line 60; figures -----	4,22,26, 39

INTERNATIONAL SEARCH REPORT

Information on patent family members

International Application No

PCT/US 00/23949

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		MX 9301917 A	31-08-1994
		WO 9320630 A	14-10-1993
		ZA 9302097 A	12-01-1994

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(19) World Intellectual Property Organization
International Bureau



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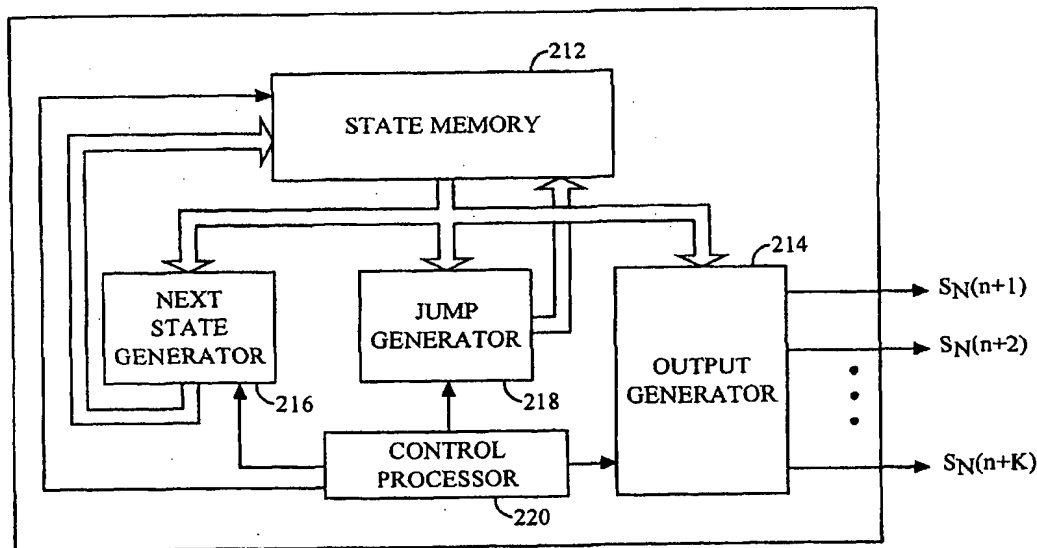
PCT

(10) International Publication Number
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(54) Title: **A METHOD AND APPARATUS FOR GENERATING MULTIPLE BITS OF A PSEUDONOISE SEQUENCE WITH EACH CLOCK PULSE BY COMPUTING THE BITS IN PARALLEL**



(57) Abstract: A novel method and apparatus for generating PN sequences with an arbitrary number of bits, where the number of bits is provided in parallel with each clock pulse is described. This allows the sequences to be generated at high speed when needed, and allows parallel processing in the acquisition and demodulation processes. In the invention, the initial values of states are loaded into registers of a parallel PN generator, which immediately generates the next n bits of the PN sequence, where n is an arbitrary number dependent on performance required. Then, a first sub-part of the PN generator (406) of the present invention receives the present state of the PN generator (406) and outputs the state of the PN generator (406) n bits in the future.

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A METHOD AND APPARATUS FOR GENERATING MULTIPLE BITS OF A PSEUDONOISE SEQUENCE WITH EACH CLOCK PULSE BY COMPUTING THE BITS IN PARALLEL

5 BACKGROUND OF THE INVENTION

I. Field of the Invention

The invention presented relates to pseudonoise (PN) sequence generators. More particularly, the present invention relates to a method and an
10 apparatus for generating PN sequence with each clock pulse by computing their bits in parallel.

II. Description of the Related Art

15 The Telecommunications Industry Association has standardized a method for code division multiple access (CDMA) communications in the IS-95 family of interim standards, entitled "Mobile Station-Base Station Compatibility Standard for Dual Mode Wideband Spread Spectrum Cellular System." In addition, the Telecommunications Industry Association in its submission to the
20 International Telecommunications Union, entitled "The cdma2000 ITU-R RTT Candidate Submission," describes proposed CDMA system that would be able to support higher data rates and higher capacity. Both in the IS-95 standard and in the cdma2000 proposal, the transmitted waveform is modulated in accordance with a pseudonoise spreading sequence.

25 The use of a pseudonoise sequence with appropriate autocorrelation characteristics is essential to the operation of a CDMA system in which multipath components are present. The generation and employment of pseudonoise sequences are described in detail in U.S. Patent No. 4,901,307, entitled "SPREAD SPECTRUM MULTIPLE ACCESS COMMUNICATION
30 SYSTEM USING SATELLITE OR TERRESTRIAL REPEATERS," assigned to the assignee of the present invention, and incorporated by reference herein. The use of CDMA techniques in a multiple access communication system is further disclosed in U.S. Patent No. 5,103,459, entitled "SYSTEM AND METHOD FOR GENERATING SIGNAL WAVEFORMS IN A CDMA CELLULAR
35 TELEPHONE SYSTEM," assigned to the assignee of the present invention, and incorporated by reference herein.

The aforementioned U.S. Patents Nos. 4,901,307 and 5,103,459 describe the use of a pilot signal used for acquisition. The use of a pilot signal enables the remote user to acquire local base station communication system in a timely manner. The remote user gets synchronization information and relative signal power information from the received pilot signal. U.S. Patents Nos. 5,644,591 and 5,805,648, both entitled "METHOD AND APPARATUS FOR PERFORMING SEARCH ACQUISITION IN A CDMA COMMUNICATION SYSTEM," describe a novel and improved method and apparatus that reduces the remote user forward link acquisition time. Both patents are assigned to the assignee of the present invention and are incorporated by reference herein.

Space or path diversity is obtained by providing multiple signal paths through simultaneous links from a remote user through two or more cell-sites. Furthermore, path diversity may be obtained by exploiting the multipath environment through spread spectrum processing by allowing a signal arriving with different propagation delays to be received and processed separately. Examples of path diversity are illustrated in U.S. Patent No. 5,101,501, entitled "SOFT HANDOFF IN A CDMA CELLULAR TELEPHONE SYSTEM," and U.S. Patent No. 5,109,390, entitled "DIVERSITY RECEIVER IN A CDMA CELLULAR TELEPHONE SYSTEM," both assigned to the assignee of the present invention, and incorporated by reference herein.

In CDMA communications systems, a pilot signal is transmitted that allows a receiver to coherently demodulate the received signal. Within demodulator of such receivers is a channel estimate generator, which estimates the channel characteristics based on the pilot signal transmitted with values known to both the transmitter and the receiver. The pilot signal is demodulated and the phase ambiguities in the received signal are resolved by taking the dot product of the received signal and the pilot signal channel estimate. An exemplary embodiment of a circuit for performing the dot product operation is disclosed in U.S. Patent No. 5,506,865, entitled "PILOT CARRIER DOT PRODUCT CIRCUIT," assigned to the assignee of the present invention, and incorporated by reference herein.

SUMMARY OF THE INVENTION

The invention presented is a novel method and apparatus for generating a PN sequences with an arbitrary number of bits, where the number of bits is provided in parallel with each clock pulse. This allows the sequences to be

generated at high speed when needed, and allows parallel processing in the acquisition and demodulation processes. The invention describes in detail generation of PN sequences as standardized for the IS-95 communications systems. As proposed in the IS-95 standards, the pseudonoise spreading sequences are maximal length sequences that are capable of being generated using linear feedback shift-registers (LSFRs). Using a linear feedback shift-register, the PN sequences are computed one bit with each clock pulse.

In the invention, the initial PN states are loaded into registers of a parallel PN generator, which immediately generates the next n bits of the PN sequence, where n is an arbitrary number dependent on performance required. In addition, the present invention provides a method of determining the register states of the parallel PN generator an arbitrary number of cycles in the future. Thus, the present invention takes the present state of the registers of the PN generator and outputs the next n bits of the generator. In addition, the PN generator of the present invention receives the present state of the PN generator and outputs the state of the PN generator n bits in the future. In this fashion, the entire PN sequence can be continuously generated.

It will be understood by one skilled in the art that although the present invention is directed toward the generation of a pseudonoise sequences compliant with systems standardized by the Telecommunications Industry Association, the teachings of the present invention are equally applicable to the generation of other pseudonoise sequences such as, the orthogonal Gold code sequences proposed for use in the W-CDMA, proposals to the International Telecommunications Industry Association, proposals by the European Telecommunications Standards Institute (ETSI), and the Association of Radio Industries and Business (ARIB).

BRIEF DESCRIPTION OF THE DRAWINGS

The features, objects, and advantages of the present invention will become more apparent from the detailed description set forth below when taken in conjunction with the drawings in which like reference characters identify correspondingly throughout and wherein:

FIG. 1 illustrates a prior art embodiment of pseudonoise (PN) generators employing linear feedback shift-registers;

FIG. 2 depicts prior art of pseudonoise generators employed to generate parallel groups of PN sequence;

FIG. 3 is a block diagram illustrating the generalized operation of the present invention apparatus for generating the PN sequences;

FIG. 4 shows one embodiment of the invention;

FIG. 5 is a simplified block diagram of an exemplary receiver chain using PN generators in accordance with the invention; and

FIG. 6 is a block diagram of a part of an exemplary single demodulation chain using PN generators in accordance with the invention.

DETAILED DESCRIPTION OF PREFERRED EMBODIMENTS

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FIG. 1a illustrates a traditional apparatus employing a linear feedback shift-register for generating pseudonoise sequences. The generalized shift-register 100 from FIG. 1a comprises memory elements 102a, 102b, . . . , 102n, holding state values $S_0(n)$, $S_1(n)$, . . . , $S_N(n)$. The last value S_N constitutes an output of the shift-register, and also a feed-back to modulo-2 adders 104a, . . . , 104m. Before the value S_N is provided to a particular modulo-2 adder 104a, . . . , 104m, it is multiplied by an associated coefficient g_0, g_1, \dots, g_N . A coefficient will take a value of '1' if a feedback is desired, and a value of '0' otherwise.

Short-code pseudonoise sequences are used to modulate and demodulate the in-phase (I) and quadrature-phase (Q) components of the CDMA waveform. The I and Q short-code PN sequences are periodic with a period of $2^{15} - 1$ with a bit stuffed at the preamble of sequence to make the sequence periodic with an even factor of 2.

The short-code PN_i sequence satisfies a linear recursion specified by the following generator polynomial (P_i):

$$P_i(x) = x^{15} + x^{13} + x^9 + x^8 + x^7 + x^5 + 1. \quad (1)$$

FIG. 1.b depicts a shift-register implementation for generating the PN_i sequence. Note that in accordance with FIG. 1a, only the '1' valued coefficients $g_{15}, g_{13}, g_9, g_8, g_7, g_5, g_0$ are present.

The short-code PN_Q sequence satisfies a linear recursion specified by the following generator polynomial (P_Q):

$$P_Q(x) = x^{15} + x^{12} + x^{11} + x^{10} + x^6 + x^5 + x^4 + x^3 + 1. \quad (2)$$

FIG. 1.c depicts a shift-register implementation for generating the PN_Q sequence.

FIG. 1c shows a shift-register implementation of a long-code PN generator with a mask. The long-code is periodic, with period $2^{42} - 1$ chips and satisfies a linear recursion specified by the following characteristic polynomial (P):

$$P(x) = x^{42} + x^{35} + x^{33} + x^{31} + x^{27} + x^{26} + x^{25} + x^{22} + x^{21} + x^{19} + x^{18} + x^{17} + x^{16} + x^{10} + x^7 + x^6 + x^5 + x^3 + x^2 + x + 1 \quad (3)$$

The mask used for the long-code is channel type dependent, and can be found along with further details about the implementation of the PN generators in a document entitled "Physical Layer Standard for cdma2000 Spread Spectrum Systems."

It is sometimes desired to obtain an output of a shift-register as a parallel combination of output state values $S_N(n)$, $S_N(n+1)$, ..., $S_N(n+K)$. FIG. 2 shows a block diagram of a parallel PN generator 200 according to the prior art. The PN generator comprises a shift-register 100 in accordance with a description for FIG. 1a, followed by a serial-to-parallel converter 202. The PN generator outputs K values of $S_N(n)$ for shift instances n , $n+1$, ..., $n+K$. However, there are K clock cycles required for generating the set of K output values. In the prior art understanding, in order to generate the parallel PN generator outputs, the outputs of the linear feedback shift-registers illustrated in FIGS. 1a and 1b are provided to the serial to parallel converter.

FIG. 3 shows a block diagram of inventive alternative to the implementation of FIG. 2. In general, a relationship between values of shift register in a state (n) and next state (n+1) can be expressed as a system of equations:

$$S_N(n+1) = g_{11} \cdot S_N(n) + \dots + g_{1N-1} \cdot S_2(n) + g_{1N} \cdot S_1(n) \quad (4a)$$

$$S_2(n+1) = g_{N-11} \cdot S_N(n) + \dots + g_{N-1N-1} \cdot S_2(n) + g_{2N} \cdot S_1(n) \quad (4n-1)$$

$$S_1(n+1) = g_{N1} \cdot S_N(n) + \dots + g_{NN-1} \cdot S_{2N-1}(n) + g_{NN} \cdot S_1(n) \quad (4n)$$

Such a system of equations can be re-written in a matrix form as:

$$S(n+1)=G*S(n), \quad (5)$$

5

where:

$S(n+1)$ is column matrix containing the state values of the state after a shift,

G is a coefficient matrix comprising the g values indicated in equations 4a-4n, and

10

$S(n)$ is a column vector of present states.

Once a state after a shift has been determined, the next state can be calculated using equation (5):

15

$$S(n+2)=G*S(n+1). \quad (6)$$

Substituting equation (5) into equation (10) then results into an equation:

$$S(n+2) = G*G*S(n) = G^2*S(n). \quad (7)$$

20

Further generalization of equation (11) yields an equation:

$$S(n+k) = G^k*S(n), \quad (8)$$

25

where k is a number expressing a state, in which an output is to be computed.

Applying these principles to FIG. 1, it is obvious that a value of a certain register in next state $S_i(n+1)$ is a function of a value of the preceding register in current state $S_{i-1}(n)$, and -- if a feedback exists -- a value of the output register in current state $S_N(n)$. Consequently, the system of equations (4) will have at most two non-zero coefficients in each of the equations (4a) through (4n).

30

As an example, the G matrix for a PN_1 shift-register in accordance with FIG. 1b will be developed as follows:

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Observing, that there is a connection between stages S_{15} and S_{14} and no feedback from stage S_{15} , it follows that the next state value of S_{15} is equal to previous state value of S_{14} . Thus, equation (4a) will take a form:

$$S_{15}(n+1) = 0 \cdot S_{15}(n) + 1 \cdot S_{14}(n) \quad (9)$$

Consequently, the first row of matrix G will contain a non-zero element only in a position g_{12} :

$$G_1 = [0100000000000000] \quad (10)$$

Equivalent relation will hold for all stages an input of which is an output of another stage.

Turning to the next stage S_{14} , one can observe that its next state value is equal to previous state value of stage S_{13} summed with a previous state value of stage S_{15} . Thus, the equation (4b) will take a form:

$$S_{14}(n+1) = 1 \cdot S_{15}(n) + 1 \cdot S_{13}(n) \quad (11)$$

Consequently, the second row of matrix G will contain a non-zero (unity) element in a position g_{21} and g_{23} :

$$G_2 = [1010000000000000] \quad (12)$$

Equivalent relation will hold between all stages an input of which is a sum of outputs of two stages.

Reference back to FIG. 3 will expand on these concepts. State memory 212 is initialized to an initial set of states $S_1(n), S_2(n), \dots, S_N(n)$. These states are then provided to an output generator 214, and a next state generator 216. Next state generator 216 contains a coefficient matrix G_{NS} formed in accordance with the principles outlined in description of equations (4) and (5). In the exemplary embodiment, the generator polynomial has relatively few feedback taps and, consequently, the resultant matrix G is sparse. This sparseness permits a relatively simple implementation of the matrix operation to be performed using fixed Boolean operator programmed into a field programmable gate array or designed into an application specific integrated circuit (ASIC).

Next state generator 216 accepts the set of states $S_1(n), S_2(n), \dots, S_N(n)$ from memory 212 to compute a set of new states $S_1(n+K), S_2(n+K), \dots, S_N(n+K)$

in accordance with equation (12), and provides the set of new states back to the state memory 212.

The output generator 214 performs a matrix operation on the current states in accordance with a matrix G_{os} formed as follows. As explained in description to FIG. 1a, the output of a shift-register is the state $S_N(n)$. From equation (8) follows that:

$$S(n+0) = G^0 S(n), \quad (13)$$

where G^0 is a matrix having non-zero elements only in the main diagonal. Inspecting the system of equations (4), it is obvious that value $S_N(n)$ can be calculated using equation (4a). This equation is equivalent to forming a row matrix G_R by taking the first row of a matrix G_{NS}^0 and multiplying it by a column matrix of states S formed from values $S_1(n), S_2(n), \dots, S_N(n)$. Therefore, the first row of a matrix G_{NS} becomes the last row of matrix G_{os} . Similarly, from equation (8), the value $S_N(n+1)$ can be calculated by forming a row matrix G_R by taking the first row of a matrix G_{NS}^2 , and multiplying it by a column matrix of states S . Thus, the last row of a matrix G_{NS} becomes the last but one row of matrix G_{os} . This process of forming the matrix G_{os} continues until all K rows are filled. In mathematical terms:

$$G_{os} = \begin{bmatrix} G_{NSL}^K \\ \vdots \\ G_{NSL}^1 \\ G_{NSL}^0 \end{bmatrix}, \quad (14)$$

where G_{NSL}^k is last row of matrix G_{NS}^k .

Once matrix G_{os} has been formed, the output generator 214 computes the values $S_N(n+1), S_N(n+2), \dots, S_N(n+K)$ by multiplying the matrix G_{os} by a column matrix of states S :

$$S_N(n+K) = G_{os} \cdot S(n) \quad (15)$$

A long-code output generator 214 differs from the structure of short-code output generator. The reason is that the long-code generator contains a mask,

which can be different for each long-code generator, see, "The cdma2000 ITU-R RTT Candidate Submission" and FIG. 1d. The PN output bit of the long code is a modulo-2 addition of values of the shift registers multiplied by the mask. The output bit can be expressed in matrix notation as follows:

5

$$pn_{out}(n) = M * S(n), \quad (16)$$

where:

$pn_{out}(n)$ is an output bit in a state n , and

10 M is a column mask matrix.

Substituting equation (8) into equation (16) results in:

$$pn_{out}(n+k) = M * G^k * S(n) \quad (17)$$

15

From equation (10) follows that desired output of $K+1$ parallel bits can be achieved by forming matrix G_{OSL}

$$G_{OSL} = \begin{bmatrix} M * G_{NSL}^K \\ \cdot \\ \cdot \\ M * G_{NSL}^1 \\ M * G_{NSL}^0 \end{bmatrix}, \quad (18)$$

20

and, once matrix G_{OSL} has been formed, the output generator 214 computes the values $pn(n)$, $pn(n+1)$, \dots , $pn(n+K)$ by multiplying the matrix G_{OSL} by a column matrix of states S :

$$pn(n+K) = G_{OSL} \cdot S(n) \quad (19)$$

25

At this point of the process the set of states, $S_1(n+K)$, $S_2(n+K)$, \dots , $S_N(n+K)$ is provided to an output generator 214, a next state generator 216, and the whole cycle is repeated.

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In particular, let us consider the G matrix for a PN_i shift-register to be the basic next state generator matrix G_{NSI} :

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$$G_{NSI} = \begin{bmatrix} 0100000000000000 \\ 1010000000000000 \\ 0001000000000000 \\ 0000100000000000 \\ 0000010000000000 \\ 1000001000000000 \\ 1000000100000000 \\ 1000000010000000 \\ 0000000001000000 \\ 1000000000100000 \\ 0000000000010000 \\ 0000000000001000 \\ 0000000000000100 \\ 0000000000000010 \\ 0000000000000001 \\ 1000000000000000 \end{bmatrix}$$

Matrix G_{NSI}^0 is as follows:

$$G_{NSI}^0 = \begin{bmatrix} 1000000000000000 \\ 0100000000000000 \\ 0010000000000000 \\ 0001000000000000 \\ 0000100000000000 \\ 0000010000000000 \\ 0000001000000000 \\ 0000000100000000 \\ 0000000010000000 \\ 0000000001000000 \\ 0000000000100000 \\ 0000000000010000 \\ 0000000000001000 \\ 0000000000000100 \\ 0000000000000010 \\ 0000000000000001 \end{bmatrix}$$

5

Taking the first row of matrix G_{NSI}^0 and last row of matrix G_{NSI} , the matrix G_{OSI} is formed as follows:

$$G_{OSI2} = \begin{bmatrix} 0100000000000000 \\ 1000000000000000 \end{bmatrix}$$

One ordinarily skilled in the art will recognize that matrix G_{OS} can be modified according to desired PN generator output, without departing from the scope of the invention. For example, if a parallel output $S_N(n)$, $S_N(n+2)$, $S_N(n+4)$, and $S_N(n+6)$ is desired, matrix G_{OS} will comprise in accordance with equation (14) first row of G_{NS}^6 in row one, first row of G_{NS}^4 in row two, first row of G_{NS}^2 in row three, and first row of G_{NS}^0 in row four.

FIG. 4 depicts a block diagram of a preferred embodiment of the parallel PN generator. In addition to the state memory 212, the output generator 214, and a next state generator 216, it contains a jump generator 218 and a control processor 220. The function of the jump generator 218 is to advance the state by predetermined number of shifts. Such a function is desirable e.g., for forward link acquisition as described in aforementioned U.S. Patent Nos. 5,644,591 and 5,805,648. In the exemplary embodiment, the PN generator is employed in a receiver in accordance to an IS-95 standard. The systems designed in accordance with an IS-95 standard comprise base stations utilizing a common PN generator, with a phase offset in increments of 64 chips for a particular pilot signal. Consequently, the jump generator 218 is functionally equivalent to next state generator 216 in that it comprises a coefficient matrix G_{JS} formed in accordance with the principles outlined in description of FIG. 1a, and raised to the power of 64.

Next state generator 216 receives the set of states $S_1(n)$, $S_2(n)$, . . . , $S_N(n)$ from memory 212 and generates a set of new states $S_1(n+64)$, $S_2(n+64)$, . . . , $S_N(n+64)$ in accordance with equation (8), and provides the set of new states back to memory 212. The reason for having a separate next state generator 216 and a jump generator 218 is that in general $K \neq L$, and, consequently, the matrices G_{OS} and G_{JS} are different. As described above, the present invention is preferably implemented in hardware adapted to the specific operation and designed to perform a specific task.

The function of the control processor 220 is to coordinate cooperation between the different subsystems, and to control bit stuffing. As described, the short-code PN sequences have a period of 2^{15} generating polynomials, and from them derived matrices, generate only sequences with period $2^{15} - 1$. The

control processor 200 monitors the output of the next state generator 216 for the state preceding the state corresponding to a period $2^{15} - 1$, for which a computation of next state according to equation (8) would exceed the state corresponding to a period $2^{15} - 1$. Once the control processor 200 detects such state it performs two operations. It will cause the output generator 214 to compute the output state values, and overwrites the last output state value with '0'. It will then avoid writing the output of the next state generator 216 into state memory 212, and will initialize the state memory 212 to initial set of states $S_1(n), S_2(n), \dots, S_N(n)$.

10

FIG. 5 depicts a simplified block diagram of an exemplary receiver chain using PN generators in accordance with the invention. The RF signal arriving at the antenna 400 is provided to the receiver (RCVR) 402, which downconverts the received signal to a baseband frequency, producing I and Q components of the signal. These components are simultaneously provided to a searcher 404 and demodulators 406a, . . . , 406c. The task of the searcher 404 is to perform searches in code space to identify candidate signals to be added to the Active Set of the remote station in order to maximize the quality of the received signal. To accomplish this task, searcher 404 will control parameters of the PN sequences generators, devised in accordance with the principles outlined in present invention. An exemplary method for performing acquisition and searching in a CDMA communication system is described in detail in aforementioned U.S. Patent Nos. 5,644,591 and 5,805,648

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In order to be effective, a receiver must be able to operate in a multipath environment and must be able to adapt to changes in physical location. In the aforementioned U.S. Patent Nos. 5,101,501 and 5,109,390, a method for exploiting the reception of multiple version of a signal is described. Demodulators 406a, 406b and 406c demodulate redundant versions of the same signal. These redundant version either correspond to multipath propagations of a signal from a single source or from multiple transmissions of the same information from multiple base stations in a soft handoff condition.

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The demodulated signals from demodulators 406a, . . . , 406c are provided to combiner 410, which combines the signals and provides them for further processing to a de-interleaver 412 and decoder 414.

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FIG. 6 illustrates the exemplary embodiment of the receiver structure of the present invention. The signal is received at antenna 400 and provided to

receiver (RCVR) 402. Receiver 402 down converts, amplifies, filters, and samples the received signal, and provides digital samples to buffer 402. In response to signals from control processor 403, a selected set of samples from buffer 404 are provided to despreader 408. In addition, in response to a signal from control processor 403, PN generator 406 provides a portion of a PN sequence to despreader 408.

Despreader 408 despreads the signal in accordance with the portion of the PN sequence provided by PN generator 406 which operates in accordance with the present invention. Within despreader 408 the PN sequence is provided to pilot despreader 412, which despreads the received signal in accordance with the portion of the short PN sequence provided by PN generator 406 and the Walsh covering sequence for the pilot signal. In the exemplary embodiment, the pilot signal is covered with the Walsh zero sequence and as such does not effect the despreading operation performed by pilot despreader 412. In addition, the portion of the short PN sequence is provided to traffic despreader 414, which despreads the signal in accordance with the short PN sequence and the Walsh traffic covering sequence W_T .

The result of the despreading operation performed by pilot despreader 412 and the result of the despreading operation performed by traffic despreader 414 are provided to dot product circuit 414. The pilot signal has known symbols and can be used to remove the phase ambiguities introduced by the propagation path as described in the aforementioned U.S. Patent No. 5,506,865. The result of the dot product operation is provided to combiner 410. Combiner 410 combines redundantly despread version of the same symbols whether transmitted by different base stations in a soft handoff environment or by the same base station traversing different propagation paths in a multipath environment.

In accordance with an exemplary demodulation chain embodiment, and previous discussion follows that a first set of matrices is required for the short-code PN generator for the I component 516, a second set for the short-code PN generator for the Q component 518, and a third set for the long-code PN generator 504.

1. Acquisition mode.

In the exemplary embodiment, the receiver is able to rapidly determine jump 64 chips ahead in the PN sequence in order to perform a correlation

14

process to determine the correlation energy of between the received signal and a portion of the PN sequence.

5 In the generation of the short PN_i sequence, state memory 212 provides the current state of the PN sequence S(n) to next state generator 216. Next state generator 216 generates the state of the PN sequence S(n+2) two cycles in advance by left-multiplying the PN sequence S(n) by the matrix G_{NSI2}:

$$G_{NSI2} = \begin{bmatrix} 10100000000000 \\ 01010000000000 \\ 00001000000000 \\ 00000100000000 \\ 10000010000000 \\ 11000001000000 \\ 11000000100000 \\ 01000000010000 \\ 10000000001000 \\ 01000000000100 \\ 000000000000100 \\ 000000000000010 \\ 000000000000001 \\ 100000000000000 \\ 010000000000000 \end{bmatrix}$$

10 In the generation of the short PN_i sequence, state memory 212 provides the current state of the PN sequence S(n) to jump generator 218. Jump generator 218 generates the state of the PN sequence S(n+2) sixty-four (64) cycles in advance by left-multiplying the PN sequence S(n) by the matrix G_{JSI64}:

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$$G_{JS164} = \begin{bmatrix} 101011010100101 \\ 010101101010010 \\ 000001100001100 \\ 000000110000110 \\ 000000011000011 \\ 000000001100001 \\ 101011010010101 \\ 011110111101111 \\ 000100001010010 \\ 000010000101001 \\ 101010010110001 \\ 110101001011000 \\ 011010100101100 \\ 101101010010110 \\ 010110101001011 \end{bmatrix}$$

In the generation of the short PN_i sequence, the next state generator 216 or the jump generator 218 provides the current state of the PN sequence $S(n)$ to output generator 214. Output generator 214 computes the values $S_N(n+1)$, $S_N(n+2)$, . . . , $S_N(n+K)$ left-multiplying a column matrix of states $S(n)$ by the matrix G_{OS12} :

$$G_{OS12} = \begin{bmatrix} 0100000000000000 \\ 1000000000000000 \end{bmatrix}$$

The short-code PN generator for the Q component 518 uses an algorithm for PN sequence generation, identical to the one for the acquisition mode. Consequently, the set of matrices as well as their application is identical.

16

 $G_{NSQ2} =$

0010000000000000
1001000000000000
1100100000000000
1100010000000000
0100001000000000
0000000100000000
0000000010000000
1000000001000000
1100000000100000
1100000000010000
1100000000001000
0100000000000100
0100000000000010
0000000000000001
1000000000000000
0100000000000000

 $G_{JSQ64} =$

1000110010111100
0100011001011110
0010001100101111
0001110100101111
1000001000101111
1100110110101111
0110011011010111
1011001101101011
1101100110110101
011000000110001
101111001000100
110100101111110
011001011100011
001100101110001
000110010111000

5

 $G_{OSQ2} = \begin{bmatrix} 0100000000000000 \\ 1000000000000000 \end{bmatrix}$

17

In the generation of the long-code PN sequence, state memory 212 provides the current state of the PN sequence $S(n)$ to next state generator 216. Next state generator 216 generates the state of the PN sequence $S(n+2)$ two cycles in advance by left-multiplying the PN sequence $S(n)$ by the matrix G_{NSL2} :

5

18

 $G_{NSL 2} =$

0010000000	0000000000	0000000000	0000000000	00
0001000000	0000000000	0000000000	0000000000	00
0000100000	0000000000	0000000000	0000000000	00
0000010000	0000000000	0000000000	0000000000	00
0000001000	0000000000	0000000000	0000000000	00
1000000100	0000000000	0000000000	0000000000	00
0100000010	0000000000	0000000000	0000000000	00
1000000001	0000000000	0000000000	0000000000	00
0100000000	1000000000	0000000000	0000000000	00
1000000000	0100000000	0000000000	0000000000	00
0100000000	0010000000	0000000000	0000000000	00
0000000000	0001000000	0000000000	0000000000	00
0000000000	0000100000	0000000000	0000000000	00
1000000000	0000010000	0000000000	0000000000	00
1100000000	0000001000	0000000000	0000000000	00
1100000000	0000000100	0000000000	0000000000	00
0100000000	0000000010	0000000000	0000000000	00
0000000000	0000000001	0000000000	0000000000	00
1000000000	0000000000	1000000000	0000000000	00
1100000000	0000000000	0100000000	0000000000	00
0100000000	0000000000	0010000000	0000000000	00
1000000000	0000000000	0001000000	0000000000	00
1100000000	0000000000	0000100000	0000000000	00
1100000000	0000000000	0000010000	0000000000	00
1100000000	0000000000	0000001000	0000000000	00
0100000000	0000000000	0000000100	0000000000	00
0000000000	0000000000	0000000010	0000000000	00
0000000000	0000000000	0000000001	0000000000	00
0000000000	0000000000	0000000000	1000000000	00
0000000000	0000000000	0000000000	0100000000	00
1000000000	0000000000	0000000000	0010000000	00
0100000000	0000000000	0000000000	0001000000	00
0000000000	0000000000	0000000000	0000100000	00
1000000000	0000000000	0000000000	0000010000	00
1100000000	0000000000	0000000000	0000001000	00
1100000000	0000000000	0000000000	0000000100	00
0100000000	0000000000	0000000000	0000000010	00
1000000000	0000000000	0000000000	0000000001	00
1100000000	0000000000	0000000000	0000000000	10
1100000000	0000000000	0000000000	0000000000	01
1100000000	0000000000	0000000000	0000000000	00
0100000000	0000000000	0000000000	0000000000	00

SUBSTITUTE SHEET (RULE 26)

In the generation of the long-code PN sequence, state memory 212 provides the current state of the PN sequence $S(n)$ to jump generator 218. Jump
5 generator 218 generates the state of the PN sequence $S(n+64)$ sixty-four (64) cycles in advance by left-multiplying the PN sequence $S(n)$ by the matrix G_{JSL64} :

20

 $G_{JSL\ 64} =$

0111001000	1101110111	1011111100	1000011111	10
1011100100	0110111011	1101111110	0100001111	11
0101110010	0011011101	1110111111	0010000111	11
0010111001	0001101110	1111011111	1001000011	11
1001011100	1000110111	0111101111	1100100001	11
1100101110	0100011011	1011110111	1110010000	11
0110010111	0010001101	1101111011	1111001000	01
0100000011	0100110001	0101000001	0111111011	10
0010000001	1010011000	1010100000	1011111101	11
0110001000	0000111011	1110101100	1101100001	01
0011000100	0000011101	1111010110	0110110000	10
1110101010	1101111001	0100010111	1011000111	11
0111010101	0110111100	1010001011	1101100011	11
1011101010	1011011110	0101000101	1110110001	11
0101110101	0101101111	0010100010	1111011000	11
0101110010	0111000000	0010101101	1111110011	11
0101110001	1110010111	1010101010	0111100110	01
0101110000	0010111100	0110101001	1011101100	10
0010111000	0001011110	0011010100	1101110110	01
1001011100	0000101111	0001101010	0110111011	00
0011100110	1101100000	0011001001	1011000010	00
0110111011	1011000111	1010011000	0101111110	10
1011011101	1101100011	1101001100	0010111111	01
0010100110	0011000110	0101011010	1001000000	00
1110011011	1100010100	1001010001	1100111111	10
0000000101	0011111101	1111010100	0110000000	01
0111001010	0100001001	0100010110	1011011111	10
0011100101	0010000100	1010001011	0101101111	11
1001110010	1001000010	0101000101	1010110111	11
1100111001	0100100001	0010100010	1101011011	11
0110011100	1010010000	1001010001	0110101101	11
0011001110	0101001000	0100101000	1011010110	11
0110101111	1111010011	1001101000	1101110100	11
1011010111	1111101001	1100110100	0110111010	01
0101101011	1111110100	1110011010	0011011101	00
1101111101	0010001101	1100110001	1001110001	00
1001110110	0100110001	0101100100	0100100111	00
0011110011	1111101111	0001001110	1010001100	00
1001111001	1111110111	1000100111	0101000110	00
1011110100	0010001100	0111101111	0010111100	10
0010110010	1100110001	1000001011	0001000001	11
1110010001	1011101111	0111111001	0000111111	01

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In the generation of the long-code PN sequence, the next state generator 216 or the jump generator 218 provides the current state of the PN sequence $S(n)$ to output generator 214. Output generator 214 first computes the output

5 state matrix G_{OSL} by left-multiplying matrix M by matrices G_{NSLO} :

22

 $G_{NSL0} =$

0100000000	0000000000	0000000000	0000000000	00
0010000000	0000000000	0000000000	0000000000	00
0001000000	0000000000	0000000000	0000000000	00
0000100000	0000000000	0000000000	0000000000	00
0000010000	0000000000	0000000000	0000000000	00
0000001000	0000000000	0000000000	0000000000	00
1000000100	0000000000	0000000000	0000000000	00
0000000010	0000000000	0000000000	0000000000	00
1000000001	0000000000	0000000000	0000000000	00
0000000000	1000000000	0000000000	0000000000	00
1000000000	0100000000	0000000000	0000000000	00
0000000000	0010000000	0000000000	0000000000	00
0000000000	0001000000	0000000000	0000000000	00
0000000000	0000100000	0000000000	0000000000	00
1000000000	0000010000	0000000000	0000000000	00
1000000000	0000001000	0000000000	0000000000	00
1000000000	0000000100	0000000000	0000000000	00
0000000000	0000000010	0000000000	0000000000	00
0000000000	0000000001	0000000000	0000000000	00
1000000000	0000000000	1000000000	0000000000	00
1000000000	0000000000	0100000000	0000000000	00
0000000000	0000000000	0010000000	0000000000	00
1000000000	0000000000	0001000000	0000000000	00
1000000000	0000000000	0000100000	0000000000	00
1000000000	0000000000	0000010000	0000000000	00
1000000000	0000000000	0000001000	0000000000	00
0000000000	0000000000	0000000100	0000000000	00
0000000000	0000000000	0000000010	0000000000	00
0000000000	0000000000	0000000001	0000000000	00
0000000000	0000000000	0000000000	1000000000	00
0000000000	0000000000	0000000000	0100000000	00
1000000000	0000000000	0000000000	0010000000	00
0000000000	0000000000	0000000000	0001000000	00
0000000000	0000000000	0000000000	0000100000	00
1000000000	0000000000	0000000000	0000010000	00
1000000000	0000000000	0000000000	0000001000	00
1000000000	0000000000	0000000000	0000000100	00
0000000000	0000000000	0000000000	0000000010	00
1000000000	0000000000	0000000000	0000000001	00
1000000000	0000000000	0000000000	0000000000	10
1000000000	0000000000	0000000000	0000000000	01
1000000000	0000000000	0000000000	0000000000	00

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and by matrix G_{NSI} :

24

 $G_{NSL1} =$

```

1000000000 0000000000 0000000000 0000000000 00
0100000000 0000000000 0000000000 0000000000 00
0010000000 0000000000 0000000000 0000000000 00
0001000000 0000000000 0000000000 0000000000 00
0000100000 0000000000 0000000000 0000000000 00
0000010000 0000000000 0000000000 0000000000 00
0000001000 0000000000 0000000000 0000000000 00
0000000100 0000000000 0000000000 0000000000 00
0000000010 0000000000 0000000000 0000000000 00
0000000001 0000000000 0000000000 0000000000 00
0000000000 1000000000 0000000000 0000000000 00
0000000000 0100000000 0000000000 0000000000 00
0000000000 0010000000 0000000000 0000000000 00
0000000000 0001000000 0000000000 0000000000 00
0000000000 0000100000 0000000000 0000000000 00
0000000000 0000010000 0000000000 0000000000 00
0000000000 0000001000 0000000000 0000000000 00
0000000000 0000000100 0000000000 0000000000 00
0000000000 0000000010 0000000000 0000000000 00
0000000000 0000000001 0000000000 0000000000 00
0000000000 0000000000 1000000000 0000000000 00
0000000000 0000000000 0100000000 0000000000 00
0000000000 0000000000 0010000000 0000000000 00
0000000000 0000000000 0001000000 0000000000 00
0000000000 0000000000 0000100000 0000000000 00
0000000000 0000000000 0000010000 0000000000 00
0000000000 0000000000 0000001000 0000000000 00
0000000000 0000000000 0000000100 0000000000 00
0000000000 0000000000 0000000010 0000000000 00
0000000000 0000000000 0000000001 0000000000 00
0000000000 0000000000 0000000000 1000000000 00
0000000000 0000000000 0000000000 0100000000 00
0000000000 0000000000 0000000000 0010000000 00
0000000000 0000000000 0000000000 0001000000 00
0000000000 0000000000 0000000000 0000100000 00
0000000000 0000000000 0000000000 0000010000 00
0000000000 0000000000 0000000000 0000001000 00
0000000000 0000000000 0000000000 0000000100 00
0000000000 0000000000 0000000000 0000000010 00
0000000000 0000000000 0000000000 0000000001 00
0000000000 0000000000 0000000000 0000000000 10
0000000000 0000000000 0000000000 0000000000 01

```

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25

, and then computes the output bits $pn_{out}(n+k)$ by multiplying the resulting matrix G_{OSL} by a column matrix of states S .

2. Demodulation mode:

5 The demodulation mode uses algorithm for PN sequence generation, identical to the one for the acquisition mode. Consequently, the set of matrices as well as their application is identical.

The short-code PN generator for the I component 516 comprises the following matrices:

10

$$G_{NS18} = \begin{bmatrix} 010010101000000 \\ 001001010100000 \\ 110110000010000 \\ 111011000001000 \\ 011101100000100 \\ 101110110000010 \\ 000101110000001 \\ 010000010000000 \\ 011010100000000 \\ 001101010000000 \\ 010100000000000 \\ 101010000000000 \\ 010101000000000 \\ 001010100000000 \\ 100101010000000 \end{bmatrix}$$

26

 $G_{JSI64} =$

101011010100101
010101101010010
000001100001100
000000110000110
000000011000011
000000001100001
101011010010101
011110111101111
000100001010010
000010000101001
101010010110001
110101001011000
011010100101100
101101010010110
010110101001011

5

 $G_{OSI8} =$

100101010000000
001010100000000
010101000000000
101010000000000
010100000000000
101000000000000
010000000000000
100000000000000

The short-code PN generator for the Q component 518 comprises the following matrices:

27

$$G_{NSQ8} = \begin{bmatrix} 101111001000000 \\ 010111100100000 \\ 101011110010000 \\ 011010110001000 \\ 000010010000100 \\ 001110000000010 \\ 100111000000001 \\ 110011100000000 \\ 111001110000000 \\ 010011110000000 \\ 000110110000000 \\ 101100010000000 \\ 111001000000000 \\ 111100100000000 \\ 011110010000000 \end{bmatrix}$$

$$G_{JSQ64} = \begin{bmatrix} 100011001011100 \\ 010001100101110 \\ 001000110010111 \\ 000111010010111 \\ 100000100010111 \\ 110011011010111 \\ 011001101101011 \\ 101100110110101 \\ 110110011011010 \\ 011000000110001 \\ 101111001000100 \\ 110100101111110 \\ 011001011100011 \\ 001100101110001 \\ 000110010111000 \end{bmatrix}$$

$$G_{os28} = \begin{bmatrix} 0111100100000000 \\ 1111001000000000 \\ 1110010000000000 \\ 1100100000000000 \\ 1001000000000000 \\ 0010000000000000 \\ 0100000000000000 \\ 1000000000000000 \end{bmatrix}$$

The long-code PN generator for 518 comprises the following matrices:

29

 $G_{NSL8} =$

0100000010	0000000000	0000000000	0000000000	00
1010000001	0000000000	0000000000	0000000000	00
0101000000	1000000000	0000000000	0000000000	00
1010100000	0100000000	0000000000	0000000000	00
0101010000	0010000000	0000000000	0000000000	00
0010101000	0001000000	0000000000	0000000000	00
1001010100	0000100000	0000000000	0000000000	00
1000101000	0000010000	0000000000	0000000000	00
0100010100	0000001000	0000000000	0000000000	00
1110001000	0000000100	0000000000	0000000000	00
1111000100	0000000010	0000000000	0000000000	00
0011100000	0000000001	0000000000	0000000000	00
1001110000	0000000000	1000000000	0000000000	00
1100111000	0000000000	0100000000	0000000000	00
1110011100	0000000000	0010000000	0000000000	00
0011001100	0000000000	0001000000	0000000000	00
0101100100	0000000000	0000100000	0000000000	00
1110110000	0000000000	0000010000	0000000000	00
1111011000	0000000000	0000001000	0000000000	00
1111101100	0000000000	0000000100	0000000000	00
1011110100	0000000000	0000000010	0000000000	00
0001111000	0000000000	0000000001	0000000000	00
1000111100	0000000000	0000000000	1000000000	00
1000011100	0000000000	0000000000	0100000000	00
0000001100	0000000000	0000000000	0010000000	00
1100000100	0000000000	0000000000	0001000000	00
0010000000	0000000000	0000000000	0000100000	00
1001000000	0000000000	0000000000	0000010000	00
1100100000	0000000000	0000000000	0000001000	00
1110010000	0000000000	0000000000	0000000100	00
0111001000	0000000000	0000000000	0000000010	00
0011100100	0000000000	0000000000	0000000001	00
1101110000	0000000000	0000000000	0000000000	10
1110111000	0000000000	0000000000	0000000000	01
0111011100	0000000000	0000000000	0000000000	00
1111101100	0000000000	0000000000	0000000000	00
1011110100	0000000000	0000000000	0000000000	00
0001111000	0000000000	0000000000	0000000000	00
1000111100	0000000000	0000000000	0000000000	00
1000011100	0000000000	0000000000	0000000000	00
1000001100	0000000000	0000000000	0000000000	00
1000000100	0000000000	0000000000	0000000000	00

30

 $G_{JSL\ 64} =$

0111001000	1101110111	1011111100	1000011111	10
1011100100	0110111011	1101111110	0100001111	11
0101110010	0011011101	1110111111	0010000111	11
0010111001	0001101110	1111011111	1001000011	11
1001011100	1000110111	0111101111	1100100001	11
1100101110	0100011011	1011110111	1110010000	11
0110010111	0010001101	1101111011	1111001000	01
0100000011	0100110001	0101000001	0111111011	10
0010000001	1010011000	1010100000	1011111101	11
0110001000	0000111011	1110101100	1101100001	01
0011000100	0000011101	1111010110	0110110000	10
1110101010	1101111001	0100010111	1011000111	11
0111010101	0110111100	1010001011	1101100011	11
1011101010	1011011110	0101000101	1110110001	11
0101110101	0101101111	0010100010	1111011000	11
0101110010	0111000000	0010101101	1111110011	11
0101110001	1110010111	1010101010	0111100110	01
0101110000	0010111100	0110101001	1011101100	10
0010111000	0001011110	0011010100	1101110110	01
1001011100	0000101111	0001101010	0110111011	00
0011100110	1101100000	0011001001	1011000010	00
0110111011	1011000111	1010011000	0101111110	10
1011011101	1101100011	1101001100	0010111111	01
0010100110	0011000110	0101011010	1001000000	00
1110011011	1100010100	1001010001	1100111111	10
0000000101	0011111101	1111010100	0110000000	01
0111001010	0100001001	0100010110	1011011111	10
0011100101	0010000100	1010001011	0101101111	11
1001110010	1001000010	0101000101	1010110111	11
1100111001	0100100001	0010100010	1101011011	11
0110011100	1010010000	1001010001	0110101101	11
0011001110	0101001000	0100101000	1011010110	11
0110101111	1111010011	1001101000	1101110100	11
1011010111	1111101001	1100110100	0110111010	01
0101101011	1111110100	1110011010	0011011101	00
1101111101	0010001101	1100110001	1001110001	00
1001110110	0100110001	0101100100	0100100111	00
0011110011	1111101111	0001001110	1010001100	00
1001111001	1111110111	1000100111	0101000110	00
1011110100	0010001100	0111101111	0010111100	10
0010110010	1100110001	1000001011	0001000001	11
1110010001	1011101111	0111111001	0000111111	01

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31

 $G_{OSL18} =$

1000000100	0000000000	0000000000	0000000000	00
0100000010	0000000000	0000000000	0000000000	00
1010000001	0000000000	0000000000	0000000000	00
0101000000	1000000000	0000000000	0000000000	00
1010100000	0100000000	0000000000	0000000000	00
0101010000	0010000000	0000000000	0000000000	00
0010101000	0001000000	0000000000	0000000000	00
0001010000	0000100000	0000000000	0000000000	00
1000101000	0000010000	0000000000	0000000000	00
1100010000	0000001000	0000000000	0000000000	00
1110001000	0000000100	0000000000	0000000000	00
0111000000	0000000010	0000000000	0000000000	00
0011100000	0000000001	0000000000	0000000000	00
1001110000	0000000000	1000000000	0000000000	00
1100111000	0000000000	0100000000	0000000000	00
0110011000	0000000000	0010000000	0000000000	00
1011001000	0000000000	0001000000	0000000000	00
1101100000	0000000000	0000100000	0000000000	00
1110110000	0000000000	0000010000	0000000000	00
1111011000	0000000000	0000001000	0000000000	00
0111101000	0000000000	0000000100	0000000000	00
0011110000	0000000000	0000000010	0000000000	00
0001111000	0000000000	0000000001	0000000000	00
0000111000	0000000000	0000000000	1000000000	00
0000011000	0000000000	0000000000	0100000000	00
1000001000	0000000000	0000000000	0010000000	00
0100000000	0000000000	0000000000	0001000000	00
0010000000	0000000000	0000000000	0000100000	00
1001000000	0000000000	0000000000	0000010000	00
1100100000	0000000000	0000000000	0000001000	00
1110010000	0000000000	0000000000	0000000100	00
0111001000	0000000000	0000000000	0000000010	00
1011100000	0000000000	0000000000	0000000001	00
1101110000	0000000000	0000000000	0000000000	10
1110111000	0000000000	0000000000	0000000000	01
1111011000	0000000000	0000000000	0000000000	00
0111101000	0000000000	0000000000	0000000000	00
0011110000	0000000000	0000000000	0000000000	00
0001111000	0000000000	0000000000	0000000000	00
0000111000	0000000000	0000000000	0000000000	00
0000011000	0000000000	0000000000	0000000000	00
0000001000	0000000000	0000000000	0000000000	00

32

 $G_{OSL\ 28} =$

0000001000	0000000000	0000000000	0000000000	00
1000000100	0000000000	0000000000	0000000000	00
0100000010	0000000000	0000000000	0000000000	00
1010000001	0000000000	0000000000	0000000000	00
0101000000	1000000000	0000000000	0000000000	00
1010100000	0100000000	0000000000	0000000000	00
0101010000	0010000000	0000000000	0000000000	00
0010100000	0001000000	0000000000	0000000000	00
0001010000	0000100000	0000000000	0000000000	00
1000100000	0000010000	0000000000	0000000000	00
1100010000	0000001000	0000000000	0000000000	00
1110000000	0000000100	0000000000	0000000000	00
0111000000	0000000010	0000000000	0000000000	00
0011100000	0000000001	0000000000	0000000000	00
1001110000	0000000000	1000000000	0000000000	00
1100110000	0000000000	0100000000	0000000000	00
0110010000	0000000000	0010000000	0000000000	00
1011000000	0000000000	0001000000	0000000000	00
1101100000	0000000000	0000100000	0000000000	00
1110110000	0000000000	0000010000	0000000000	00
1111010000	0000000000	0000001000	0000000000	00
0111100000	0000000000	0000000100	0000000000	00
0011110000	0000000000	0000000010	0000000000	00
0001110000	0000000000	0000000001	0000000000	00
0000110000	0000000000	0000000000	1000000000	00
0000010000	0000000000	0000000000	0100000000	00
1000000000	0000000000	0000000000	0010000000	00
0100000000	0000000000	0000000000	0001000000	00
0010000000	0000000000	0000000000	0000100000	00
1001000000	0000000000	0000000000	0000010000	00
1100100000	0000000000	0000000000	0000001000	00
1110010000	0000000000	0000000000	0000000100	00
0111000000	0000000000	0000000000	0000000010	00
1011100000	0000000000	0000000000	0000000001	00
1101110000	0000000000	0000000000	0000000000	10
1110110000	0000000000	0000000000	0000000000	01
1111010000	0000000000	0000000000	0000000000	00
0111100000	0000000000	0000000000	0000000000	00
0011110000	0000000000	0000000000	0000000000	00
0001110000	0000000000	0000000000	0000000000	00
0000110000	0000000000	0000000000	0000000000	00
0000010000	0000000000	0000000000	0000000000	00

33

 $G_{OSL\ 38} =$

0000010000	0000000000	0000000000	0000000000	00
0000001000	0000000000	0000000000	0000000000	00
1000000100	0000000000	0000000000	0000000000	00
0100000010	0000000000	0000000000	0000000000	00
1010000001	0000000000	0000000000	0000000000	00
0101000000	1000000000	0000000000	0000000000	00
1010100000	0100000000	0000000000	0000000000	00
0101000000	0010000000	0000000000	0000000000	00
0010100000	0001000000	0000000000	0000000000	00
0001000000	0000100000	0000000000	0000000000	00
1000100000	0000010000	0000000000	0000000000	00
1100000000	0000001000	0000000000	0000000000	00
1110000000	0000000100	0000000000	0000000000	00
0111000000	0000000010	0000000000	0000000000	00
0011100000	0000000001	0000000000	0000000000	00
1001100000	0000000000	1000000000	0000000000	00
1100100000	0000000000	0100000000	0000000000	00
0110000000	0000000000	0010000000	0000000000	00
1011000000	0000000000	0001000000	0000000000	00
1101100000	0000000000	0000100000	0000000000	00
1110100000	0000000000	0000010000	0000000000	00
1111000000	0000000000	0000001000	0000000000	00
0111100000	0000000000	0000000100	0000000000	00
0011100000	0000000000	0000000010	0000000000	00
0001100000	0000000000	0000000001	0000000000	00
0000100000	0000000000	0000000000	1000000000	00
0000000000	0000000000	0000000000	0100000000	00
1000000000	0000000000	0000000000	0010000000	00
0100000000	0000000000	0000000000	0001000000	00
0010000000	0000000000	0000000000	0000100000	00
1001000000	0000000000	0000000000	0000010000	00
1100100000	0000000000	0000000000	0000001000	00
1110000000	0000000000	0000000000	0000000100	00
0111000000	0000000000	0000000000	0000000010	00
1011100000	0000000000	0000000000	0000000001	00
1101100000	0000000000	0000000000	0000000000	10
1110100000	0000000000	0000000000	0000000000	01
1111000000	0000000000	0000000000	0000000000	00
0111100000	0000000000	0000000000	0000000000	00
0011100000	0000000000	0000000000	0000000000	00
0001100000	0000000000	0000000000	0000000000	00
0000100000	0000000000	0000000000	0000000000	00

34

 $G_{OSL\ 48} =$

0000100000	0000000000	0000000000	0000000000	00
0000010000	0000000000	0000000000	0000000000	00
0000001000	0000000000	0000000000	0000000000	00
1000000100	0000000000	0000000000	0000000000	00
0100000010	0000000000	0000000000	0000000000	00
1010000001	0000000000	0000000000	0000000000	00
0101000000	1000000000	0000000000	0000000000	00
1010000000	0100000000	0000000000	0000000000	00
0101000000	0010000000	0000000000	0000000000	00
0010000000	0001000000	0000000000	0000000000	00
0001000000	0000100000	0000000000	0000000000	00
1000000000	0000010000	0000000000	0000000000	00
1100000000	0000001000	0000000000	0000000000	00
1110000000	0000000100	0000000000	0000000000	00
0111000000	0000000010	0000000000	0000000000	00
0011000000	0000000001	0000000000	0000000000	00
1001000000	0000000000	1000000000	0000000000	00
1100000000	0000000000	0100000000	0000000000	00
0110000000	0000000000	0010000000	0000000000	00
1011000000	0000000000	0001000000	0000000000	00
1101000000	0000000000	0000100000	0000000000	00
1110000000	0000000000	0000010000	0000000000	00
0111000000	0000000000	0000001000	0000000000	00
0011000000	0000000000	0000000100	0000000000	00
0001000000	0000000000	0000000010	0000000000	00
0000100000	0000000000	0000000001	0000000000	00
0000010000	0000000000	0000000000	1000000000	00
0000001000	0000000000	0000000000	0100000000	00
1000000000	0000000000	0000000000	0010000000	00
0100000000	0000000000	0000000000	0001000000	00
0010000000	0000000000	0000000000	0000100000	00
1001000000	0000000000	0000000000	0000010000	00
1100000000	0000000000	0000000000	0000001000	00
1110000000	0000000000	0000000000	0000000100	00
0111000000	0000000000	0000000000	0000000010	00
1011000000	0000000000	0000000000	0000000001	00
1101000000	0000000000	0000000000	0000000000	10
1110000000	0000000000	0000000000	0000000000	01
1111000000	0000000000	0000000000	0000000000	00
0111000000	0000000000	0000000000	0000000000	00
0011000000	0000000000	0000000000	0000000000	00
0001000000	0000000000	0000000000	0000000000	00

35

 $G_{OSL\ 58} =$

0001000000	0000000000	0000000000	0000000000	00
0000100000	0000000000	0000000000	0000000000	00
0000010000	0000000000	0000000000	0000000000	00
0000001000	0000000000	0000000000	0000000000	00
1000000100	0000000000	0000000000	0000000000	00
0100000010	0000000000	0000000000	0000000000	00
1010000001	0000000000	0000000000	0000000000	00
0100000000	1000000000	0000000000	0000000000	00
1010000000	0100000000	0000000000	0000000000	00
0100000000	0010000000	0000000000	0000000000	00
0010000000	0001000000	0000000000	0000000000	00
0000000000	0000100000	0000000000	0000000000	00
1000000000	0000010000	0000000000	0000000000	00
1100000000	0000001000	0000000000	0000000000	00
1110000000	0000000100	0000000000	0000000000	00
0110000000	0000000010	0000000000	0000000000	00
0010000000	0000000001	0000000000	0000000000	00
1000000000	0000000000	1000000000	0000000000	00
1100000000	0000000000	0100000000	0000000000	00
0110000000	0000000000	0010000000	0000000000	00
1010000000	0000000000	0001000000	0000000000	00
1100000000	0000000000	0000100000	0000000000	00
1110000000	0000000000	0000010000	0000000000	00
1110000000	0000000000	0000001000	0000000000	00
0110000000	0000000000	0000000100	0000000000	00
0010000000	0000000000	0000000010	0000000000	00
0000000000	0000000000	0000000001	0000000000	00
0000000000	0000000000	0000000000	1000000000	00
0000000000	0000000000	0000000000	0100000000	00
1000000000	0000000000	0000000000	0010000000	00
0100000000	0000000000	0000000000	0001000000	00
0010000000	0000000000	0000000000	0000100000	00
1000000000	0000000000	0000000000	0000010000	00
1100000000	0000000000	0000000000	0000001000	00
1110000000	0000000000	0000000000	0000000100	00
0110000000	0000000000	0000000000	0000000010	00
1010000000	0000000000	0000000000	0000000001	00
1100000000	0000000000	0000000000	0000000000	10
1110000000	0000000000	0000000000	0000000000	01
1110000000	0000000000	0000000000	0000000000	00
0110000000	0000000000	0000000000	0000000000	00
0010000000	0000000000	0000000000	0000000000	00

36

 $G_{OSL} =$

001000000	000000000	000000000	000000000	00
000100000	000000000	000000000	000000000	00
000010000	000000000	000000000	000000000	00
000001000	000000000	000000000	000000000	00
000000100	000000000	000000000	000000000	00
100000100	000000000	000000000	000000000	00
010000010	000000000	000000000	000000000	00
100000001	000000000	000000000	000000000	00
010000000	100000000	000000000	000000000	00
100000000	010000000	000000000	000000000	00
010000000	001000000	000000000	000000000	00
000000000	000100000	000000000	000000000	00
000000000	000010000	000000000	000000000	00
100000000	000001000	000000000	000000000	00
110000000	000000100	000000000	000000000	00
110000000	000000010	000000000	000000000	00
010000000	000000001	000000000	000000000	00
000000000	000000000	100000000	000000000	00
100000000	000000000	010000000	000000000	00
010000000	000000000	001000000	000000000	00
100000000	000000000	000100000	000000000	00
110000000	000000000	000010000	000000000	00
110000000	000000000	000001000	000000000	00
110000000	000000000	000000100	000000000	00
010000000	000000000	000000010	000000000	00
000000000	000000000	000000001	000000000	00
000000000	000000000	000000000	100000000	00
000000000	000000000	000000000	010000000	00
100000000	000000000	000000000	001000000	00
010000000	000000000	000000000	000100000	00
000000000	000000000	000000000	000010000	00
100000000	000000000	000000000	000001000	00
110000000	000000000	000000000	000000100	00
110000000	000000000	000000000	000000010	00
010000000	000000000	000000000	000000001	00
100000000	000000000	000000000	000000000	10
110000000	000000000	000000000	000000000	01
110000000	000000000	000000000	000000000	00
010000000	000000000	000000000	000000000	00

37

 $G_{OSL\ 78} =$

0100000000	0000000000	0000000000	0000000000	00
0010000000	0000000000	0000000000	0000000000	00
0001000000	0000000000	0000000000	0000000000	00
0000100000	0000000000	0000000000	0000000000	00
0000010000	0000000000	0000000000	0000000000	00
0000001000	0000000000	0000000000	0000000000	00
1000000100	0000000000	0000000000	0000000000	00
0000000010	0000000000	0000000000	0000000000	00
1000000001	0000000000	0000000000	0000000000	00
0000000000	1000000000	0000000000	0000000000	00
1000000000	0100000000	0000000000	0000000000	00
0000000000	0010000000	0000000000	0000000000	00
0000000000	0001000000	0000000000	0000000000	00
0000000000	0000100000	0000000000	0000000000	00
1000000000	0000010000	0000000000	0000000000	00
1000000000	0000001000	0000000000	0000000000	00
1000000000	0000000100	0000000000	0000000000	00
0000000000	0000000010	0000000000	0000000000	00
0000000000	0000000001	0000000000	0000000000	00
1000000000	0000000000	1000000000	0000000000	00
1000000000	0000000000	0100000000	0000000000	00
0000000000	0000000000	0010000000	0000000000	00
1000000000	0000000000	0001000000	0000000000	00
1000000000	0000000000	0000100000	0000000000	00
1000000000	0000000000	0000010000	0000000000	00
1000000000	0000000000	0000001000	0000000000	00
0000000000	0000000000	0000000100	0000000000	00
0000000000	0000000000	0000000010	0000000000	00
0000000000	0000000000	0000000001	0000000000	00
0000000000	0000000000	0000000000	1000000000	00
0000000000	0000000000	0000000000	0100000000	00
1000000000	0000000000	0000000000	0010000000	00
0000000000	0000000000	0000000000	0001000000	00
0000000000	0000000000	0000000000	0000100000	00
1000000000	0000000000	0000000000	0000010000	00
1000000000	0000000000	0000000000	0000001000	00
1000000000	0000000000	0000000000	0000000100	00
0000000000	0000000000	0000000000	0000000010	00
1000000000	0000000000	0000000000	0000000001	00
1000000000	0000000000	0000000000	0000000000	10
1000000000	0000000000	0000000000	0000000000	01
1000000000	0000000000	0000000000	0000000000	00

38

 $G_{OSL} =$

1000000000	0000000000	0000000000	0000000000	00
0100000000	0000000000	0000000000	0000000000	00
0010000000	0000000000	0000000000	0000000000	00
0001000000	0000000000	0000000000	0000000000	00
0000100000	0000000000	0000000000	0000000000	00
0000010000	0000000000	0000000000	0000000000	00
0000001000	0000000000	0000000000	0000000000	00
0000000100	0000000000	0000000000	0000000000	00
0000000010	0000000000	0000000000	0000000000	00
0000000001	0000000000	0000000000	0000000000	00
0000000000	1000000000	0000000000	0000000000	00
0000000000	0100000000	0000000000	0000000000	00
0000000000	0010000000	0000000000	0000000000	00
0000000000	0001000000	0000000000	0000000000	00
0000000000	0000100000	0000000000	0000000000	00
0000000000	0000010000	0000000000	0000000000	00
0000000000	0000001000	0000000000	0000000000	00
0000000000	0000000100	0000000000	0000000000	00
0000000000	0000000010	0000000000	0000000000	00
0000000000	0000000001	0000000000	0000000000	00
0000000000	0000000000	1000000000	0000000000	00
0000000000	0000000000	0100000000	0000000000	00
0000000000	0000000000	0010000000	0000000000	00
0000000000	0000000000	0001000000	0000000000	00
0000000000	0000000000	0000100000	0000000000	00
0000000000	0000000000	0000010000	0000000000	00
0000000000	0000000000	0000001000	0000000000	00
0000000000	0000000000	0000000100	0000000000	00
0000000000	0000000000	0000000010	0000000000	00
0000000000	0000000000	0000000001	0000000000	00
0000000000	0000000000	0000000000	1000000000	00
0000000000	0000000000	0000000000	0100000000	00
0000000000	0000000000	0000000000	0010000000	00
0000000000	0000000000	0000000000	0001000000	00
0000000000	0000000000	0000000000	0000100000	00
0000000000	0000000000	0000000000	0000010000	00
0000000000	0000000000	0000000000	0000001000	00
0000000000	0000000000	0000000000	0000000100	00
0000000000	0000000000	0000000000	0000000010	00
0000000000	0000000000	0000000000	0000000001	00
0000000000	0000000000	0000000000	0000000000	10
0000000000	0000000000	0000000000	0000000000	01

The previous description of the preferred embodiments is provided to enable any person skilled in the art to make or use the present invention. The

various modifications to these embodiments will be readily apparent to those skilled in the art, and the generic principles defined herein may be applied to other embodiments without the use of the inventive faculty. Thus, the present invention is not intended to be limited to the embodiments shown herein but is
5 to be accorded the widest scope consistent with the principles and novel features disclosed herein.

WHAT IS CLAIMED IS:

1. An apparatus for generating multiple bits of a pseudonoise sequence with each clock pulse by computing the bits in parallel, comprising:
 - a) a state memory;
 - b) a next state generator communicatively connected with said state memory; and
 - c) an output generator communicatively connected with said state memory and said next state generator.
2. The apparatus of claim 1 wherein said state memory has been configured to hold:
 - a) a set of initial values of states; and
 - b) a set of values of states generated by said next state generator or a jump generator.
3. The apparatus of claim 1 wherein said set of initial values of states comprises:
 - a) coefficients of a generating polynomial.
4. The apparatus of claim 3 wherein said generating polynomial is:

$$P_1(x) = x^{15} + x^{13} + x^9 + x^8 + x^7 + x^5 + 1$$
5. The apparatus of claim 3 wherein said generating polynomial is:

$$P_0(x) = x^{15} + x^{12} + x^{11} + x^{10} + x^6 + x^5 + x^4 + x^3 + 1$$
6. The apparatus of claim 3 wherein said generating polynomial is:

$$P(x) = x^{42} + x^{35} + x^{33} + x^{31} + x^{27} + x^{26} + x^{25} + x^{22} + x^{21} + x^{19} + x^{18} + x^{17} + x^{16} + (1) + x^{10} + x^7 + x^6 + x^5 + x^3 + x^2 + x + 1.$$

41

2 7. The apparatus of claim 1 wherein said next state generator has been configured to:

- 4 a) accept one set of values of states;
 b) generate another set of values of states a first pre-determined
 6 number of clocks apart from current state by multiplying said accepted values by a next step matrix; and
 8 c) provide said another set of values of states to said memory and said output generator.

2 8. The apparatus of claim 7 wherein said first pre-determined number of clocks is two and said next step matrix G_{NS12} is:

$$G_{NS12} = \begin{bmatrix} 10100000000000 \\ 01010000000000 \\ 00001000000000 \\ 00000100000000 \\ 10000010000000 \\ 11000001000000 \\ 11000000100000 \\ 01000000010000 \\ 10000000001000 \\ 01000000000100 \\ 00000000000010 \\ 00000000000001 \\ 10000000000000 \\ 01000000000000 \end{bmatrix}$$

4

2 9. The apparatus of claim 7 wherein said first pre-determined number of clocks is two and said next step matrix G_{NSQ2} is:

$G_{NSQ2} =$

00100000000000
10010000000000
11001000000000
11000100000000
01000010000000
00000001000000
00000000100000
10000000010000
11000000001000
11000000000100
11000000000010
01000000000001
00000000000000
10000000000000
01000000000000

4

- 2 10. The apparatus of claim 7 wherein said a first pre-determined number of
clocks is eight and said next step matrix G_{NS18} is:

4

 $G_{NS18} =$

010010101000000
001001010100000
110110000010000
111011000001000
011101100000100
101110110000010
000101110000001
010000010000000
011010100000000
001101010000000
010100000000000
101010000000000
010101000000000
001010100000000
100101010000000

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- 2 11. The apparatus of claim 7 wherein said a first pre-determined number of
clocks is eight and said next step matrix G_{NSQ2} is:

$$G_{NSQ8} = \begin{bmatrix} 101111001000000 \\ 010111100100000 \\ 101011110010000 \\ 011010110001000 \\ 000010010000100 \\ 001110000000010 \\ 100111000000001 \\ 110011100000000 \\ 111001110000000 \\ 010011110000000 \\ 000110110000000 \\ 101100010000000 \\ 111001000000000 \\ 111100100000000 \\ 011110010000000 \end{bmatrix}$$

4

- 2 12. The apparatus of claim 1 wherein said output generator has been
configured to:

- 4 a) one set of values of states; and
b) generate multiple output bits in parallel by multiplying said
6 accepted values by an output state matrix.

- 2 13. The apparatus of claim 12 wherein said multiple is two and said output
state matrix G_{OS12} is:

4

$$G_{OS12} = \begin{bmatrix} 010000000000000 \\ 100000000000000 \end{bmatrix}$$

- 2 14. The apparatus of claim 12 wherein said multiple is two and said output
state matrix G_{OSQ2} is:

4

$$G_{OSQ2} = \begin{bmatrix} 010000000000000 \\ 100000000000000 \end{bmatrix}$$

- 2 15. The apparatus of claim 12 wherein said multiple is eight and said output
state matrix G_{OS18} is:

4

$$G_{OS18} = \begin{bmatrix} 1001010100000000 \\ 0010101000000000 \\ 0101010000000000 \\ 1010100000000000 \\ 0101000000000000 \\ 1010000000000000 \\ 0100000000000000 \\ 1000000000000000 \end{bmatrix}$$

- 2 16. The apparatus of claim 12 wherein said multiple is eight and said output
state matrix G_{OS08} is:

4

$$G_{OS08} = \begin{bmatrix} 0111100100000000 \\ 1111001000000000 \\ 1110010000000000 \\ 1100100000000000 \\ 1001000000000000 \\ 0010000000000000 \\ 0100000000000000 \\ 1000000000000000 \end{bmatrix}$$

- 2 17. The apparatus of claim 1 further comprising a jump generator.
- 2 18. The apparatus of claim 17 wherein said jump generator has been
configured to:
- 4 a) accept one set of values of states;
- 4 b) generate values of states a second pre-determined number of
6 clocks apart from current state by multiplying said accepted values by a jump
state matrix; and

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- 8 c) provide said values of states to said memory and said output generator.

- 2 19. The apparatus of claim 18 wherein said second pre-determined number is sixty-four and said jump state matrix G_{JS164} is:

$$G_{JS164} = \begin{bmatrix} 101011010100101 \\ 010101101010010 \\ 000001100001100 \\ 000000110000110 \\ 000000011000011 \\ 000000001100001 \\ 101011010010101 \\ 011110111101111 \\ 000100001010010 \\ 000010000101001 \\ 101010010110001 \\ 110101001011000 \\ 011010100101100 \\ 101101010010110 \\ 010110101001011 \end{bmatrix}$$

4

20. The apparatus of claim 18 wherein said second pre-determined number is sixty-four and said jump state matrix G_{JSQ64} is:

2

$$G_{JSQ64} = \begin{bmatrix} 100011001011100 \\ 010001100101110 \\ 001000110010111 \\ 000111010010111 \\ 100000100010111 \\ 110011011010111 \\ 011001101101011 \\ 101100110110101 \\ 110110011011010 \\ 011000000110001 \\ 101111001000100 \\ 110100101111110 \\ 011001011100011 \\ 001100101110001 \\ 000110010111000 \end{bmatrix}$$

21. The apparatus of claim 1 further comprising a controller.
22. The apparatus of claim 21 wherein said controller has been configured to
 2 monitor output bits of said next state generator for a pre-determined
 combination, and when said pre-determined combination has been reached to:
 - 4 a) overwrite an appropriate output bit value with a value of '0';
 - b) void writing values of states generated by said next state
 6 generator to said state memory; and
 - c) instruct said state memory to provide a set of initial values of
 8 states to said next state generator.
23. A pseudonoise (PN) sequence generator comprising:
 - 2 a) state memory for storing at least one state of a PN generator
 polynomial;
 - 4 b) next state generator for receiving said at least one state of said PN
 generator polynomial and for generating a second state of said PN generator
 polynomial by performing a matrix operation upon said at least one state of
 6 said PN generator polynomial; and
 - 8 c) output generator for receiving said at least one state of said PN
 generator polynomial and for generating at least one PN sequence output by

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10 performing a matrix operation upon said at least one state of said PN generator polynomial.

24. The apparatus of Claim 23 wherein said at least one state comprises the
2 fifteen component state of a PN short code.

25. The apparatus of Claim 23 wherein said at least one state comprises the
2 forty two component state of a PN long code.

26. The apparatus of Claim 23 wherein said generator polynomial (P_I) is:
2 $P_I(x) = x^{15} + x^{13} + x^9 + x^8 + x^7 + x^5 + 1$

27. The apparatus of Claim 23 wherein said generator polynomial (P_Q) is:
2 $P_Q(x) = x^{15} + x^{12} + x^{11} + x^{10} + x^6 + x^5 + x^4 + x^3 + 1$

28. The apparatus of Claim 23 wherein said generator polynomial (P) is:
2 $P(x) = x^{42} + x^{35} + x^{33} + x^{31} + x^{27} + x^{26} + x^{25} + x^{22} + x^{21} + x^{19} + x^{18} + x^{17} + x^{16} +$
 $(1) + x^{10} + x^7 + x^6 + x^5 + x^3 + x^2 + x + 1.$

29. The apparatus of Claim 23 wherein said next state generator computes
2 the state of PN sequence generator two clock cycles in the future and performs
said matrix operation in accordance with the matrix GNSI21:

48

$$G_{NSI2} = \begin{bmatrix} 10100000000000 \\ 01010000000000 \\ 00001000000000 \\ 00000100000000 \\ 10000010000000 \\ 11000001000000 \\ 11000000100000 \\ 01000000010000 \\ 10000000001000 \\ 01000000000100 \\ 000000000000100 \\ 000000000000010 \\ 000000000000001 \\ 100000000000000 \\ 010000000000000 \end{bmatrix}$$

4

30. The apparatus of Claim 23 wherein said next state generator performs
2 said matrix operation in accordance with the matrix GNSQ2:

$$G_{NSQ2} = \begin{bmatrix} 00100000000000 \\ 10010000000000 \\ 11001000000000 \\ 11000100000000 \\ 01000010000000 \\ 00000001000000 \\ 00000000100000 \\ 10000000010000 \\ 11000000001000 \\ 11000000000100 \\ 110000000000100 \\ 010000000000010 \\ 000000000000001 \\ 100000000000000 \\ 010000000000000 \end{bmatrix}$$

31. The apparatus of Claim 23 wherein said output generator computes the next two outputs of said PN sequence generator and performs said matrix operation in accordance with the matrix GOSI2:

$$G_{OSI2} = \begin{bmatrix} 0100000000000000 \\ 1000000000000000 \end{bmatrix}.$$

32. The apparatus of Claim 23 wherein said PN generator programmed into an ASIC.

33. The apparatus of Claim 23 wherein said PN generator programmed into a field programmable gate array.

34. A method for generating multiple bits of a pseudonoise sequence with each clock pulse by computing the bits in parallel, comprising the steps of:

- a) storing at least one set of values of states in a state memory;
- b) generating a second set of values of states by a next state generator, said second set being derived from said at least one set; and
- c) generating a set of output bits in parallel by an output generator, said set of output bits being derived from said at least one set of values of states.

35. The method of claim 34, wherein the step of storing at least one set of values of states comprises the steps of:

- a) holding a set of initial values of states; and
- b) holding another set of values of states from said next state generator or from a jump generator.

36. The method of claim 34, wherein the step of generating a second set of values of states comprises the step of:

- a) multiplying said at least one set of values of states by a next step matrix.

37. The method of claim 34, wherein the step of generating a set of output bits in parallel comprises the step of:

50

- 4 a) multiplying said at least one set of values of states by an output state matrix.

2 38. The method of claim 34, further comprising the step of monitoring a set of values of states of said next state generator for a pre-determined combination.

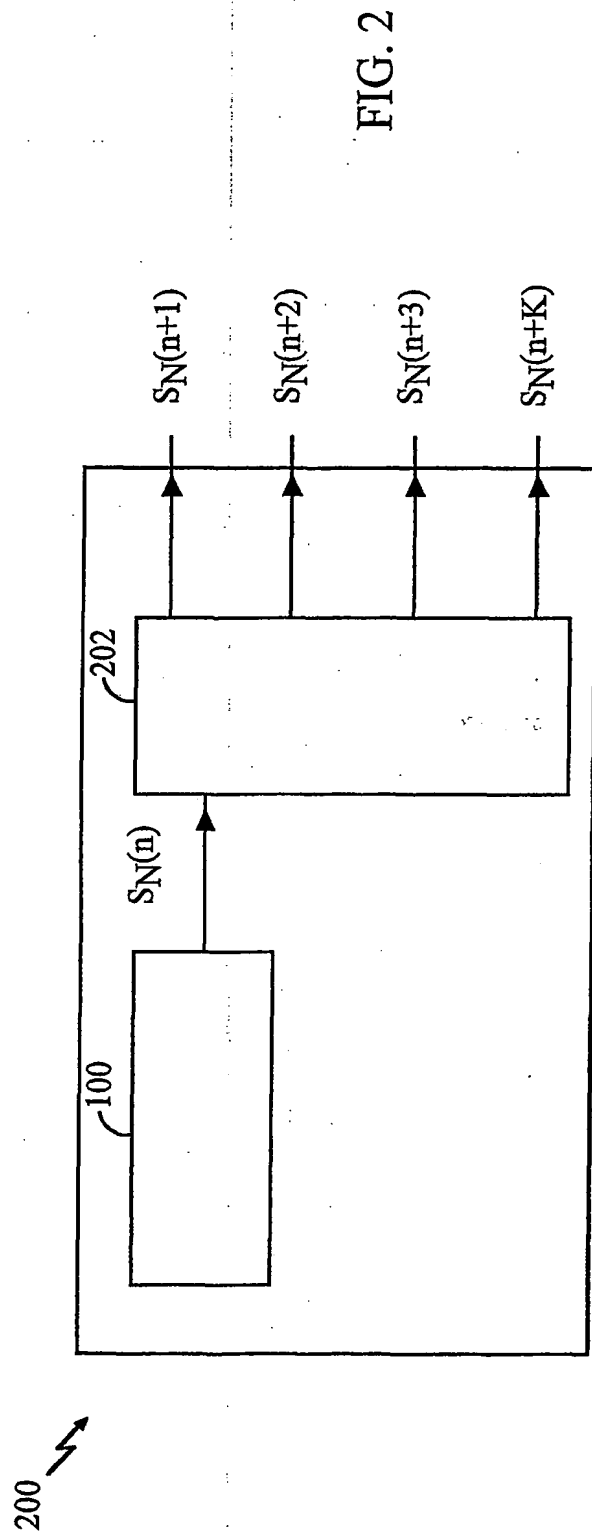
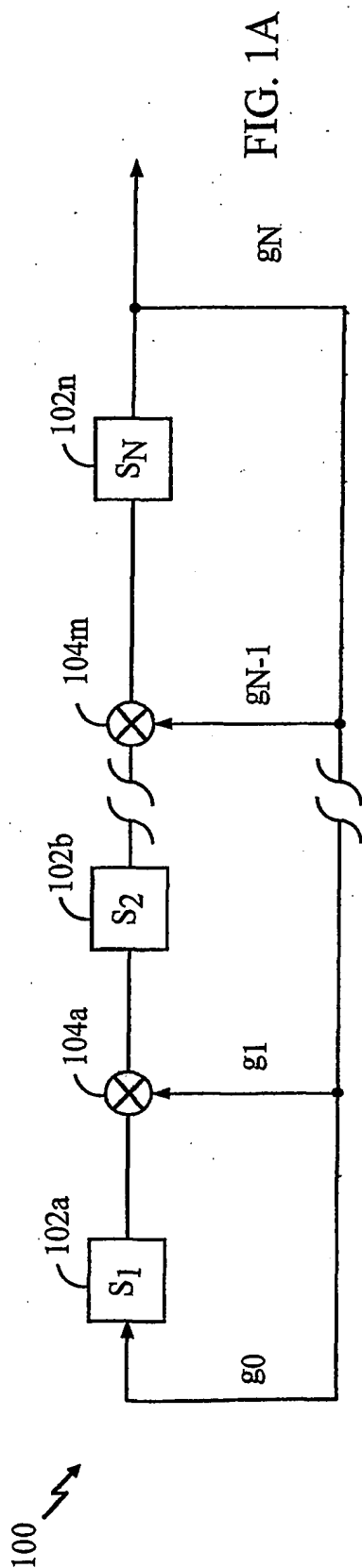
2 39. The method of claim 38, wherein upon detecting said pre-determined combination, the method further comprises the steps of:

- 4 a) overwriting an appropriate output bit value with a value of '0';
4 b) voiding writing said second set of values of states generated by said next state generator to said state memory; and
6 c) instructing said state memory to provide a set of initial values of states to said next state generator.

2 40. The method of claim 34, further comprising the step of generating a third set of values of states by a jump state generator, said second set being derived from said at least one set.

2 41. The method of claim 40, wherein the step of generating a third set of values of states by a jump state generator comprises the step of:

- 4 a) multiplying said at least one set of values of states by a jump state matrix.



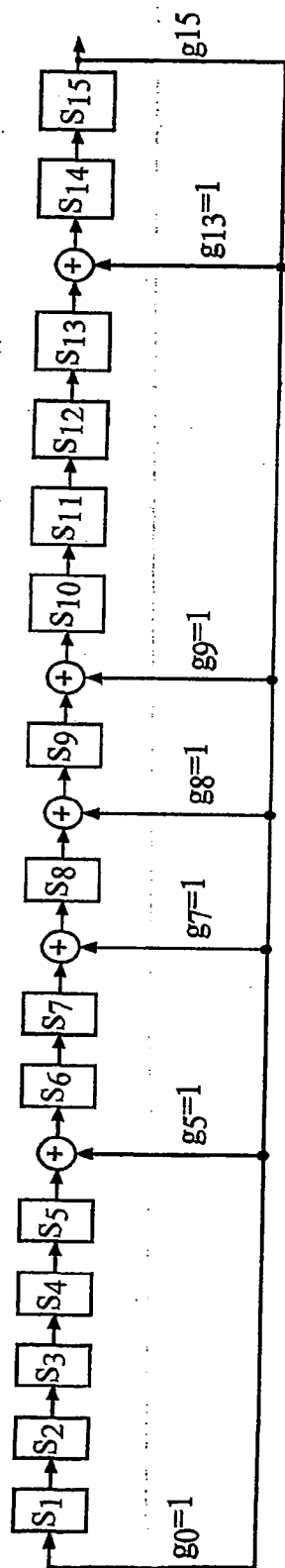


FIG. 1B

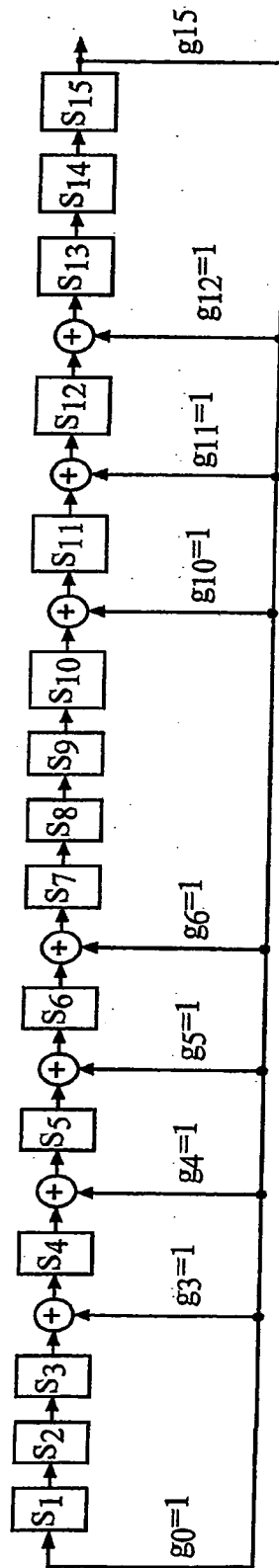


FIG. 1C

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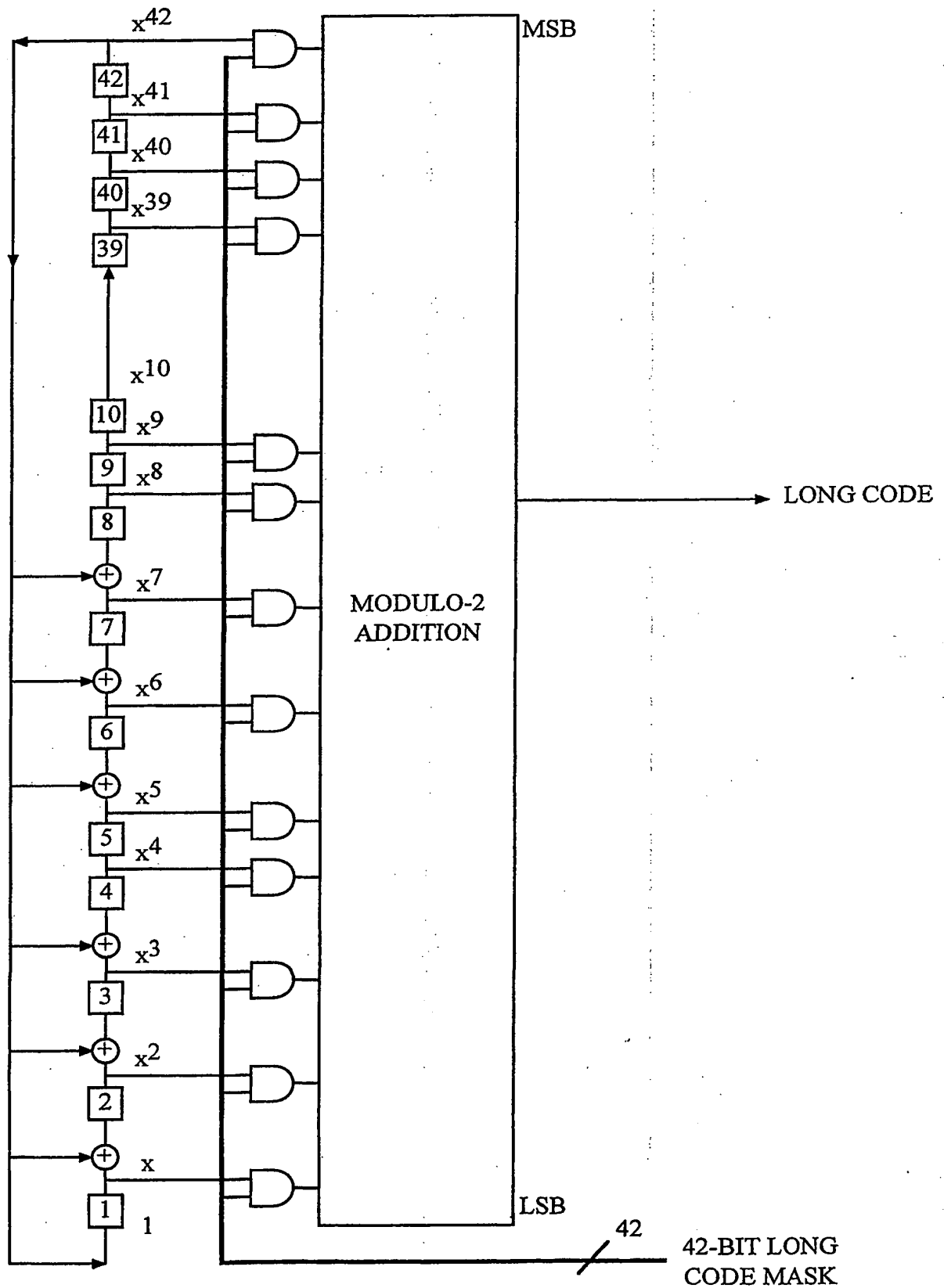


FIG. 1D

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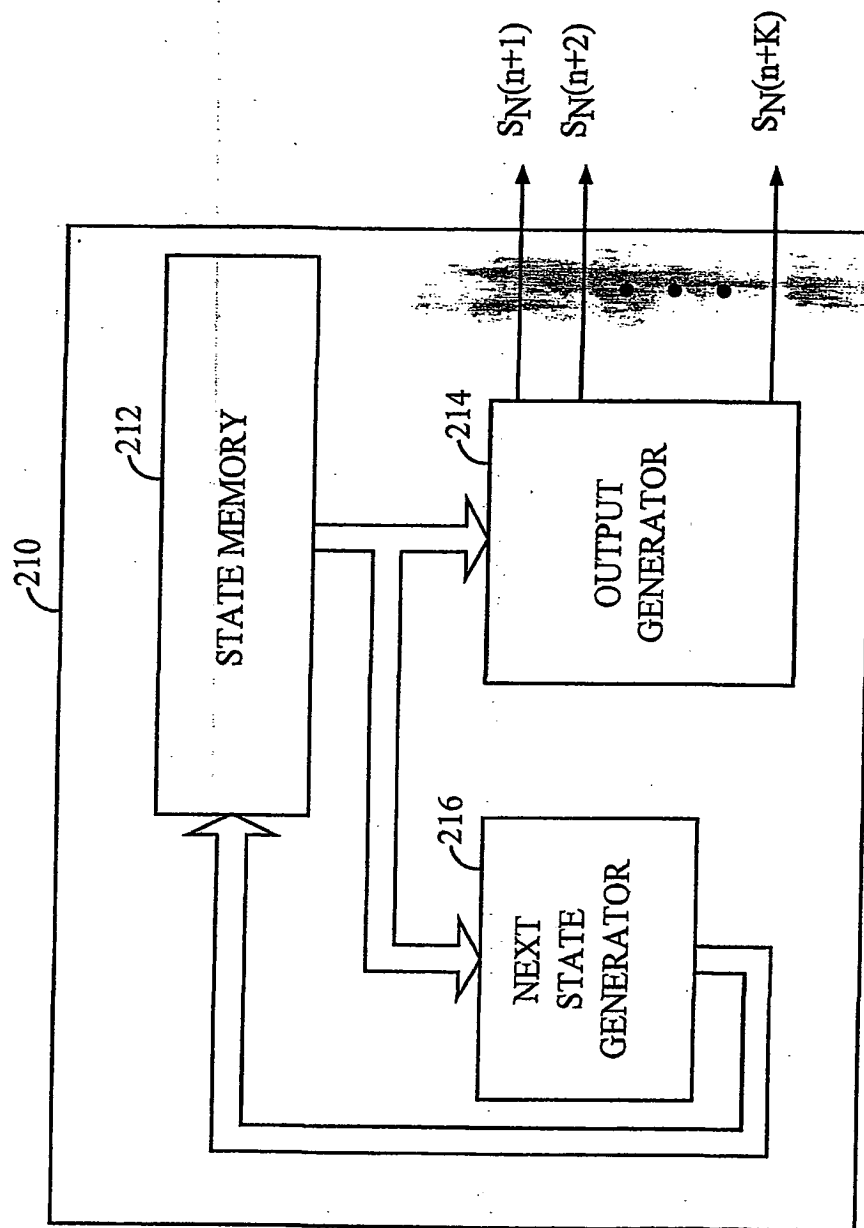


FIG. 3

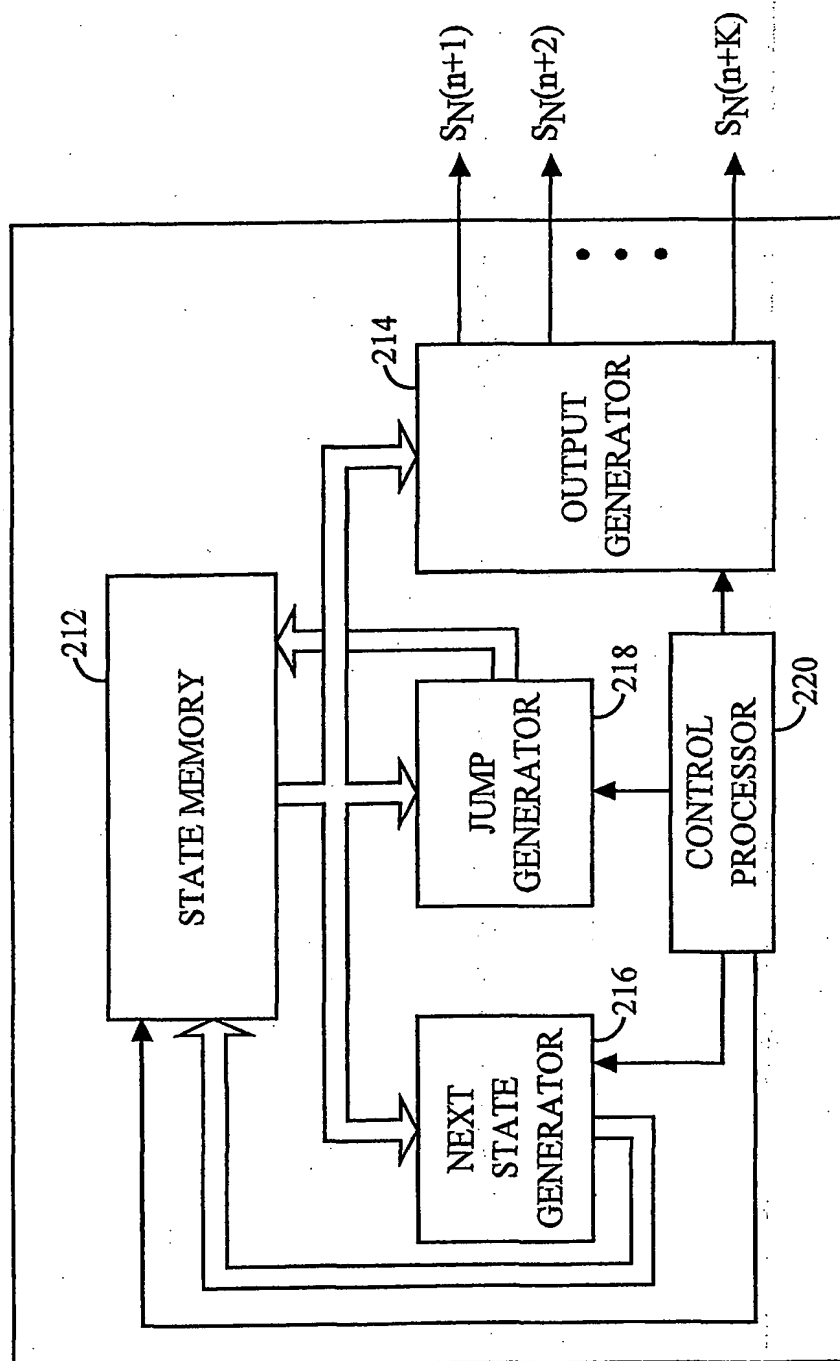


FIG. 4

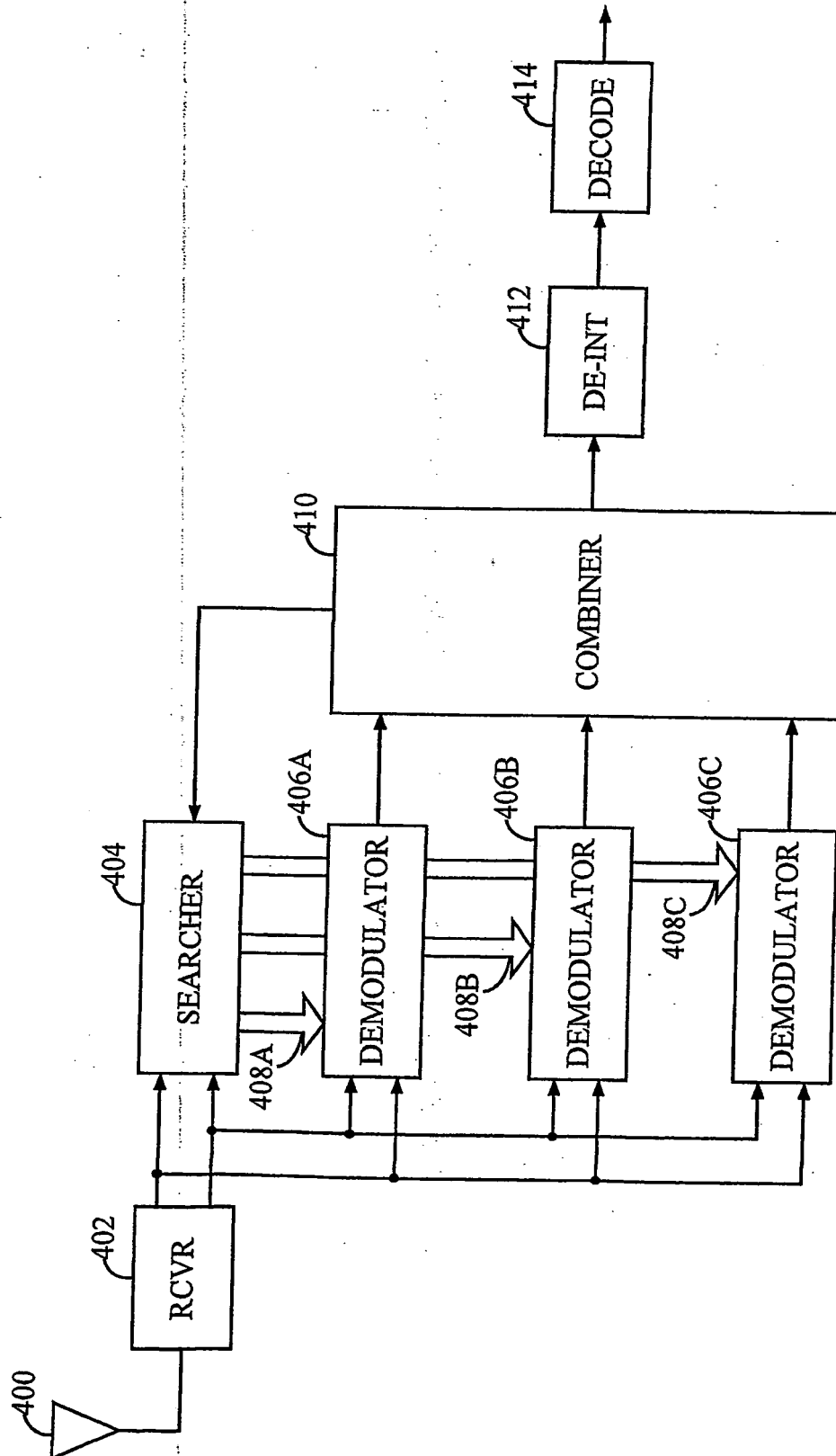


FIG. 5

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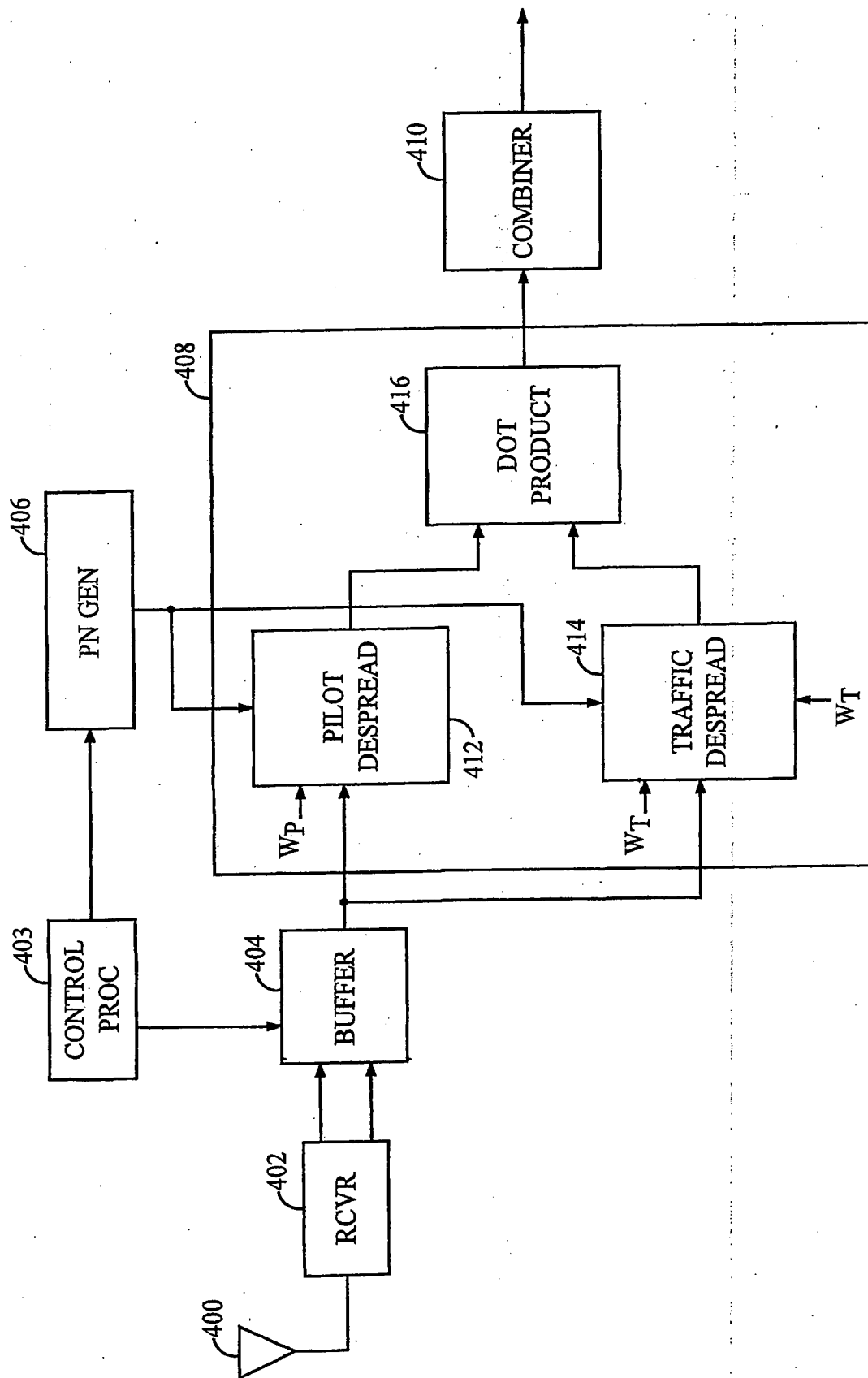


FIG. 6

PCT/US 00/23949

INTERNATIONAL SEARCH REPORT

International Application No

PCT/US 00/23949

C.(Continuation) DOCUMENTS CONSIDERED TO BE RELEVANT

Category *	Citation of document, with indication, where appropriate, of the relevant passages	Relevant to claim No.
A	US 5 228 054 A (GILHOUSEN KLEIN S ET AL) 13 July 1993 (1993-07-13) column 4, line 37 - line 60; figures -----	4, 22, 26, 39

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